



# RDF Pattern Matching using Sortable Views

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# Outline

- RDF Pattern Rewriting using Views
- Sortable View
- Rewriting using Sortable Views
- Evaluation



# RDF Pattern Rewriting using View

subject	predicate	object
Alice	Marry	Bob
Alice	isMother	Chris
Chris	coAuthor	Bob
Chris	coAuthor	Frank
Alice	isMother	David
...	....	...



# RDF Pattern Rewriting using View

subject	predicate	object
Alice	Marry	Bob
Alice	isMother	Chris
Chris	coAuthor	Bob
Chris	coAuthor	Frank
Alice	isMother	David
...	....	...

Fathers and sons who are coauthored with each other?



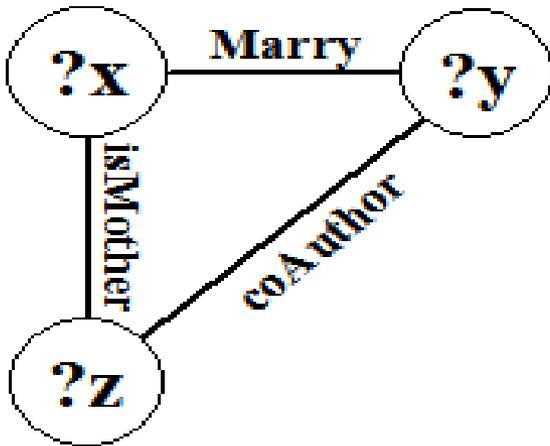
# RDF Pattern Rewriting using View

subject	predicate	object
Alice	Marry	Bob
Alice	isMother	Chris
Chris	coAuthor	Bob
Chris	coAuthor	Frank
Alice	isMother	David
...	....	...

$(?x, \text{Marry}, ?y), (?x, \text{isMother}, ?z), (?y, \text{coAuthor}, ?z)$



# RDF Pattern Rewriting using View

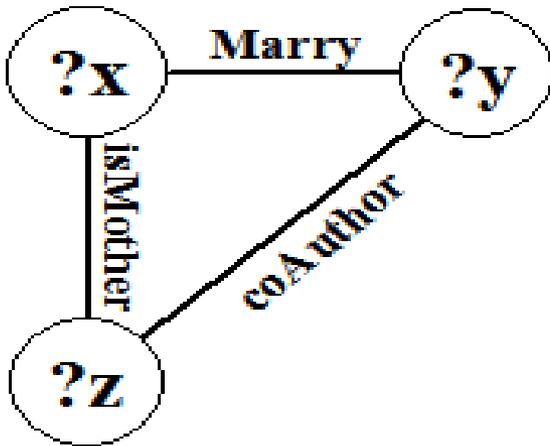


**View  $V_1$**

$(?x, \text{Marry}, ?y), (?x, \text{isMother}, ?z), (?y, \text{coAuthor}, ?z)$



# RDF Pattern Rewriting using View



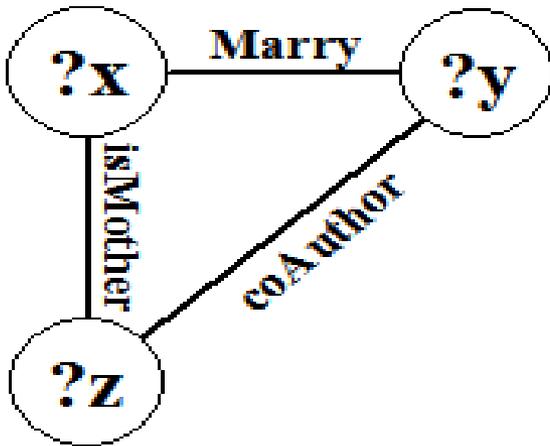
**View  $V_1$**

Husbands and their wives  
(?x, Marry, ?y)

(?x, Marry, ?y), (?x, isMother, ?z), (?y, coAuthor, ?z)



# RDF Pattern Rewriting using View



**View  $V_1$**

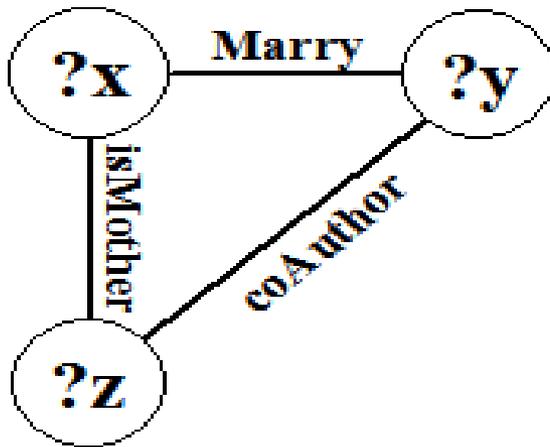
Husbands and their wives  
( $?x$ , Marry,  $?y$ )

Mothers and their sons  
( $?x$ , isMother,  $?z$ )

( $?x$ , Marry,  $?y$ ), ( $?x$ , isMother,  $?z$ ), ( $?y$ , coAuthor,  $?z$ )



# RDF Pattern Rewriting using View



**View  $V_1$**

Join on  $?x$ :  
Fathers ( $?y$ ) and their sons ( $?z$ )

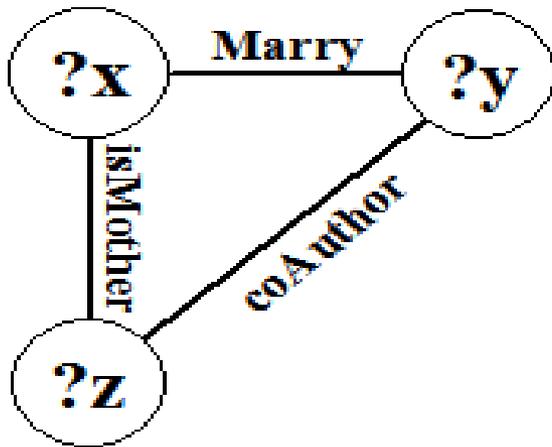
Husbands and their wives  
( $?x$ , Marry,  $?y$ )

Mothers and their sons  
( $?x$ , isMother,  $?z$ )

$(?x, Marry, ?y), (?x, isMother, ?z), (?y, coAuthor, ?z)$



# RDF Pattern Rewriting using View



**View  $V_1$**

Join on  $?x$ :  
Fathers ( $?y$ ) and their sons ( $?z$ )

Coauthors  
( $?z$ , coAuthor,  $?y$ )

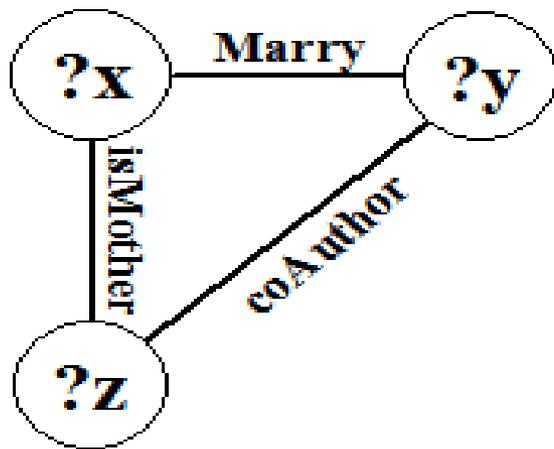
Husbands and their wives  
( $?x$ , Marry,  $?y$ )

Mothers and their sons  
( $?x$ , isMother,  $?z$ )

$(?x, Marry, ?y), (?x, isMother, ?z), (?y, coAuthor, ?z)$



# RDF Pattern Rewriting using View



**View  $V_1$**

Join on  $?y$  and  $?z$ :  
Coauthored fathers and sons

Join on  $?x$ :  
Fathers ( $?y$ ) and their sons ( $?z$ )

Coauthors  
( $?z$ , coAuthor,  $?y$ )

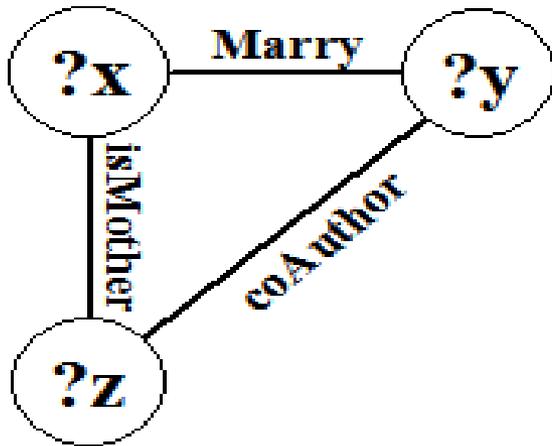
Husbands and their wives  
( $?x$ , Marry,  $?y$ )

Mothers and their sons  
( $?x$ , isMother,  $?z$ )

$(?x, Marry, ?y), (?x, isMother, ?z), (?y, coAuthor, ?z)$



# RDF Pattern Rewriting using View



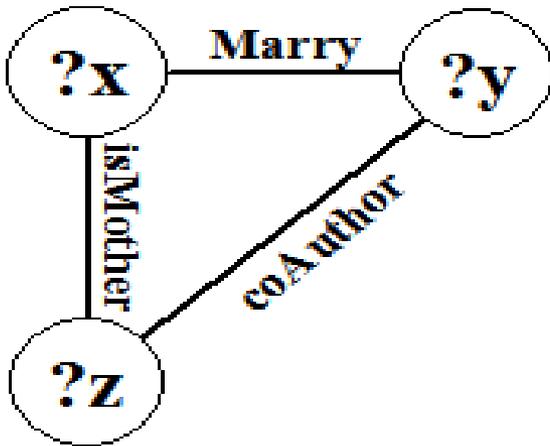
**View  $V_1$**

$[V_1]$	$?x$	$?y$	$?z$
	Alice	Bob	Chris
	...	...	...

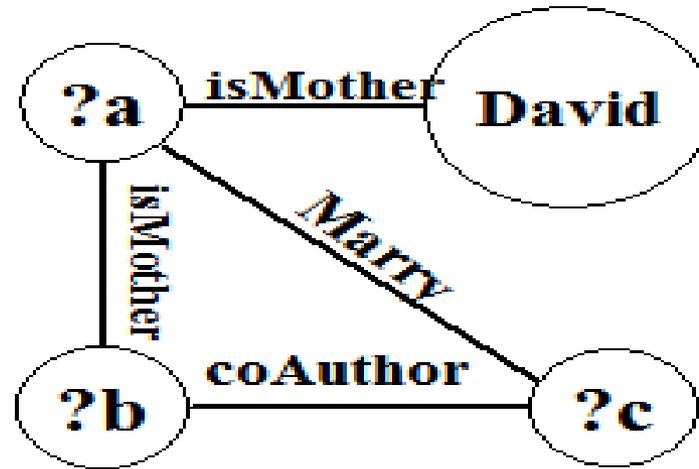
$(?x, \text{Marry}, ?y), (?x, \text{isMother}, ?z), (?y, \text{coAuthor}, ?z)$



# RDF Pattern Rewriting using View



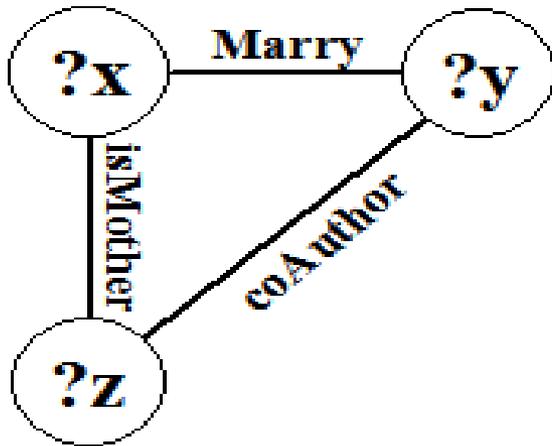
**View  $V_1$**



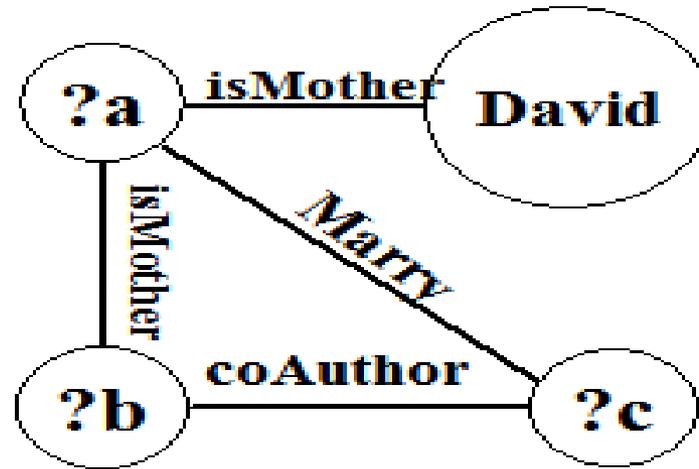
**Pattern  $Q_1$**



# RDF Pattern Rewriting using View



**View  $V_1$**

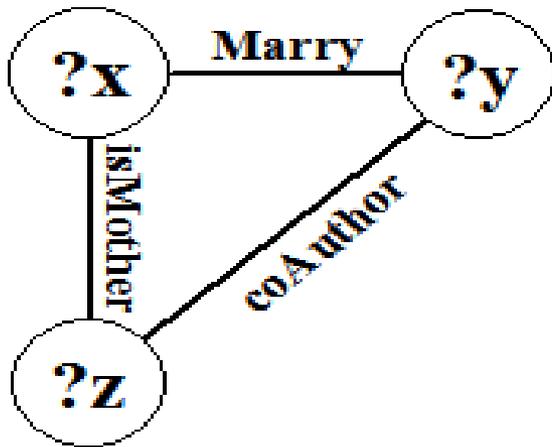


**Pattern  $Q_1$**

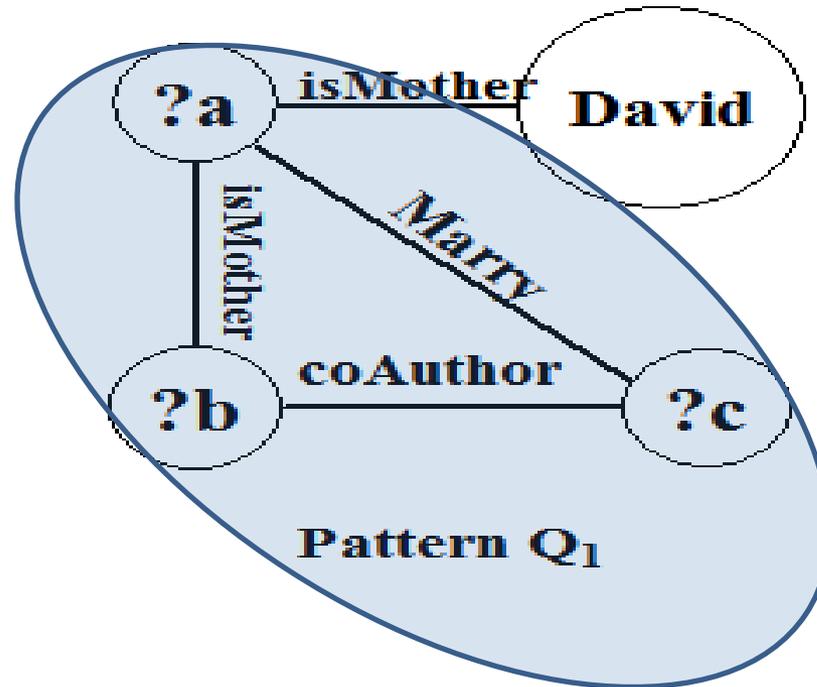
David's brothers who are coauthored with their fathers?



# RDF Pattern Rewriting using View



**View  $V_1$**

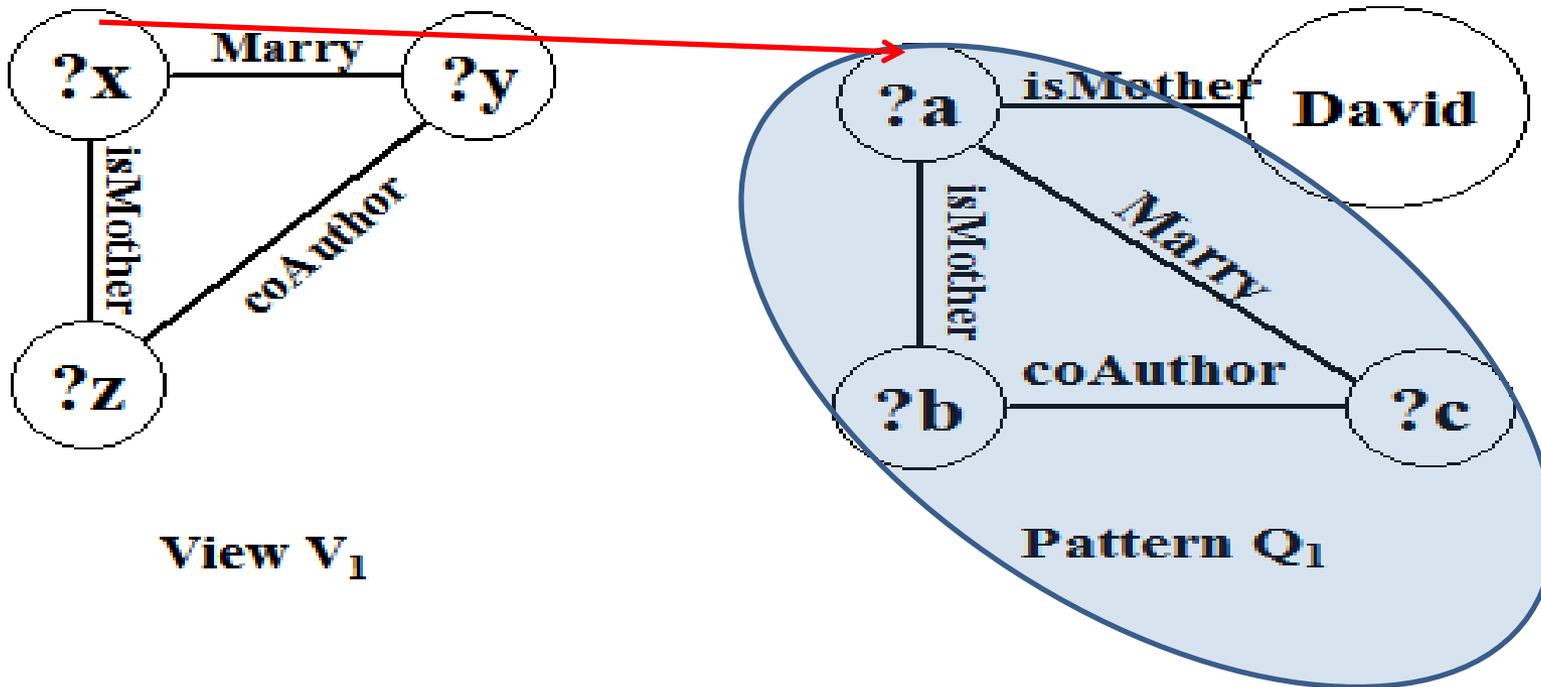


**Pattern  $Q_1$**

David's brothers who are coauthored with their fathers?



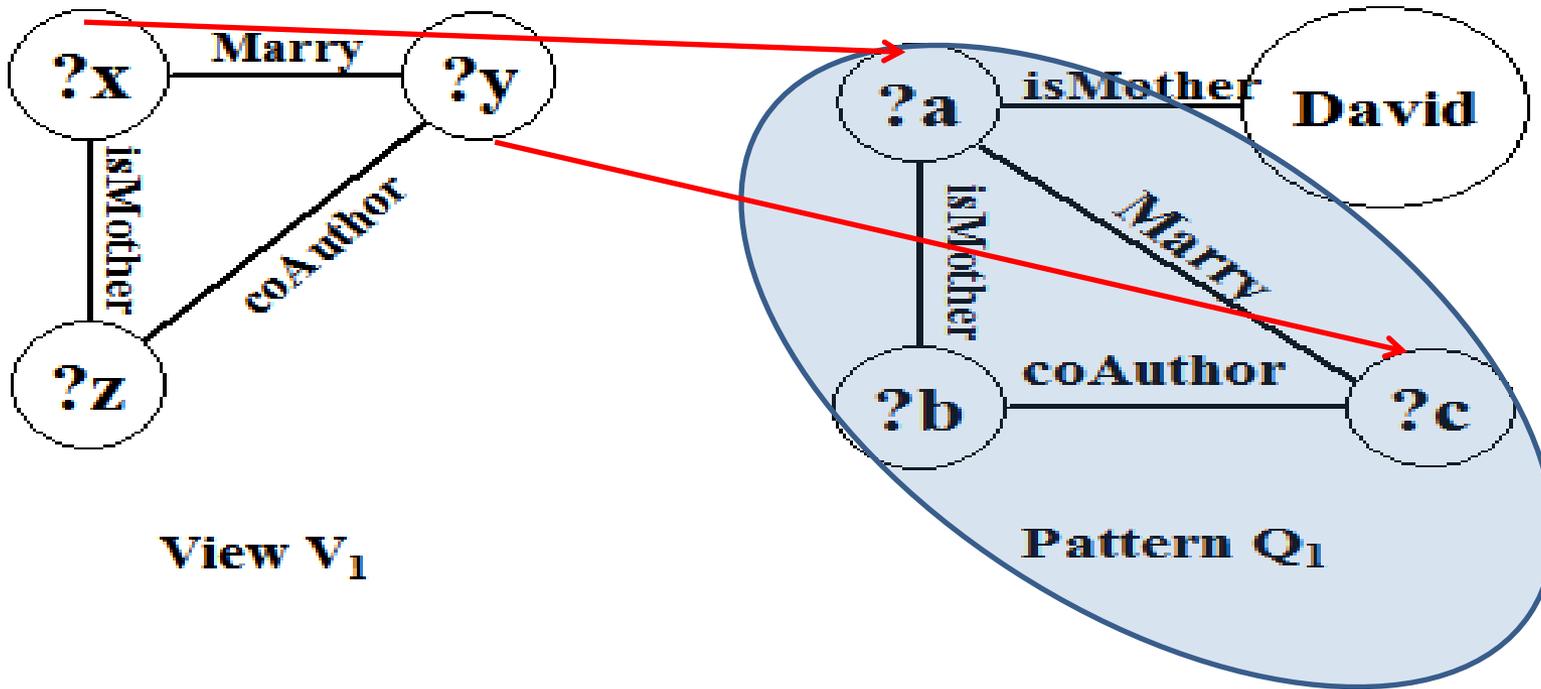
# RDF Pattern Rewriting using View



David's brothers who are coauthored with their fathers?



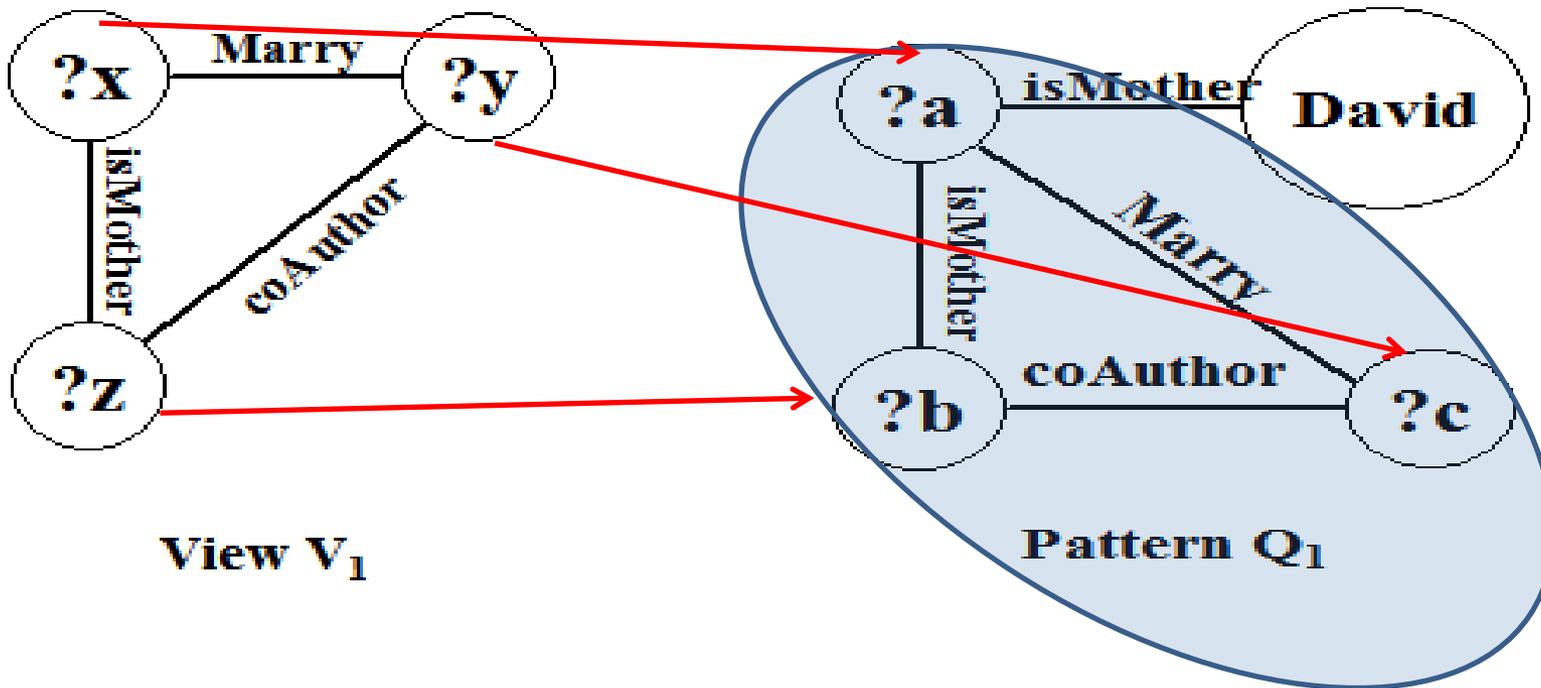
# RDF Pattern Rewriting using View



David's brothers who are coauthored with their fathers?



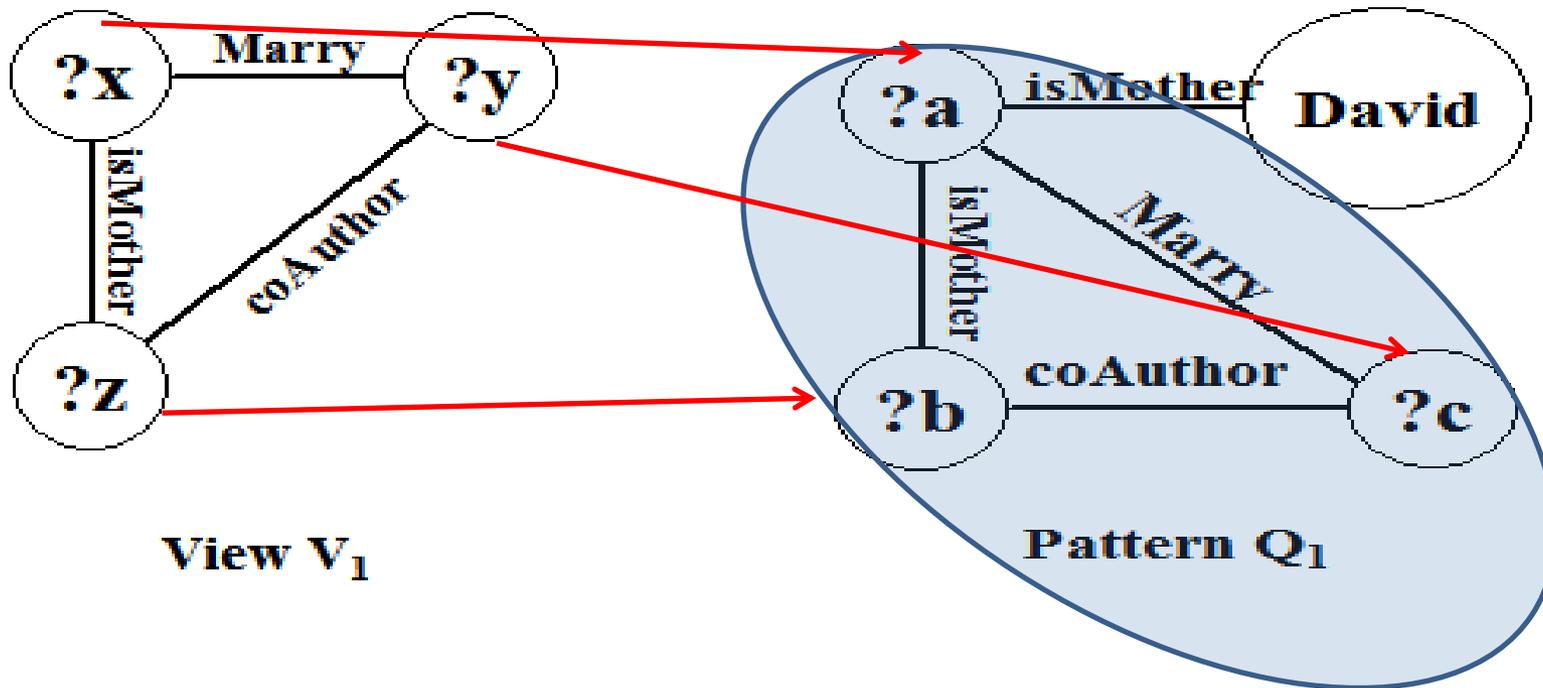
# RDF Pattern Rewriting using View



David's brothers who are coauthored with their fathers?



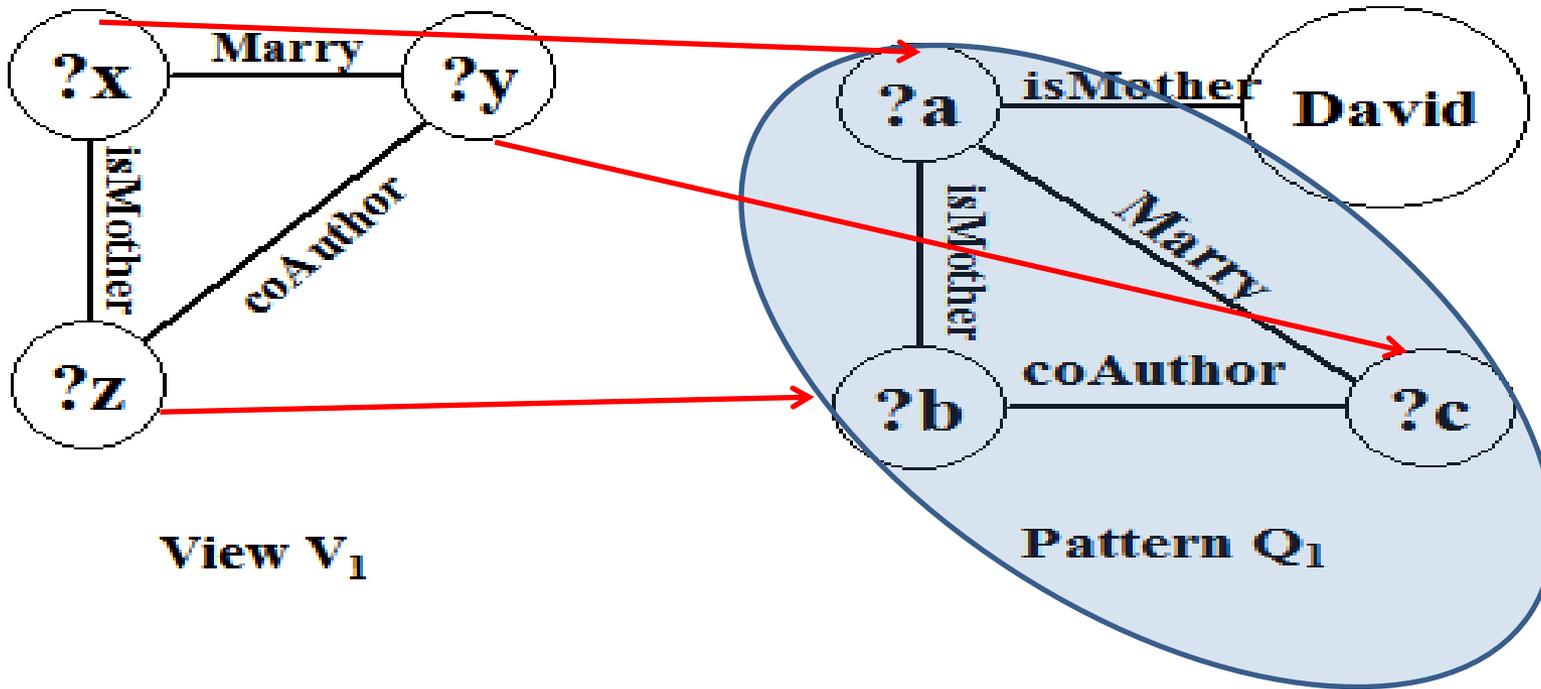
# RDF Pattern Rewriting using View



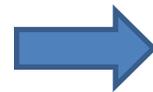
V1	?x	?y	?z
	Alice	Bob	Chris
	...	...	...



# RDF Pattern Rewriting using View



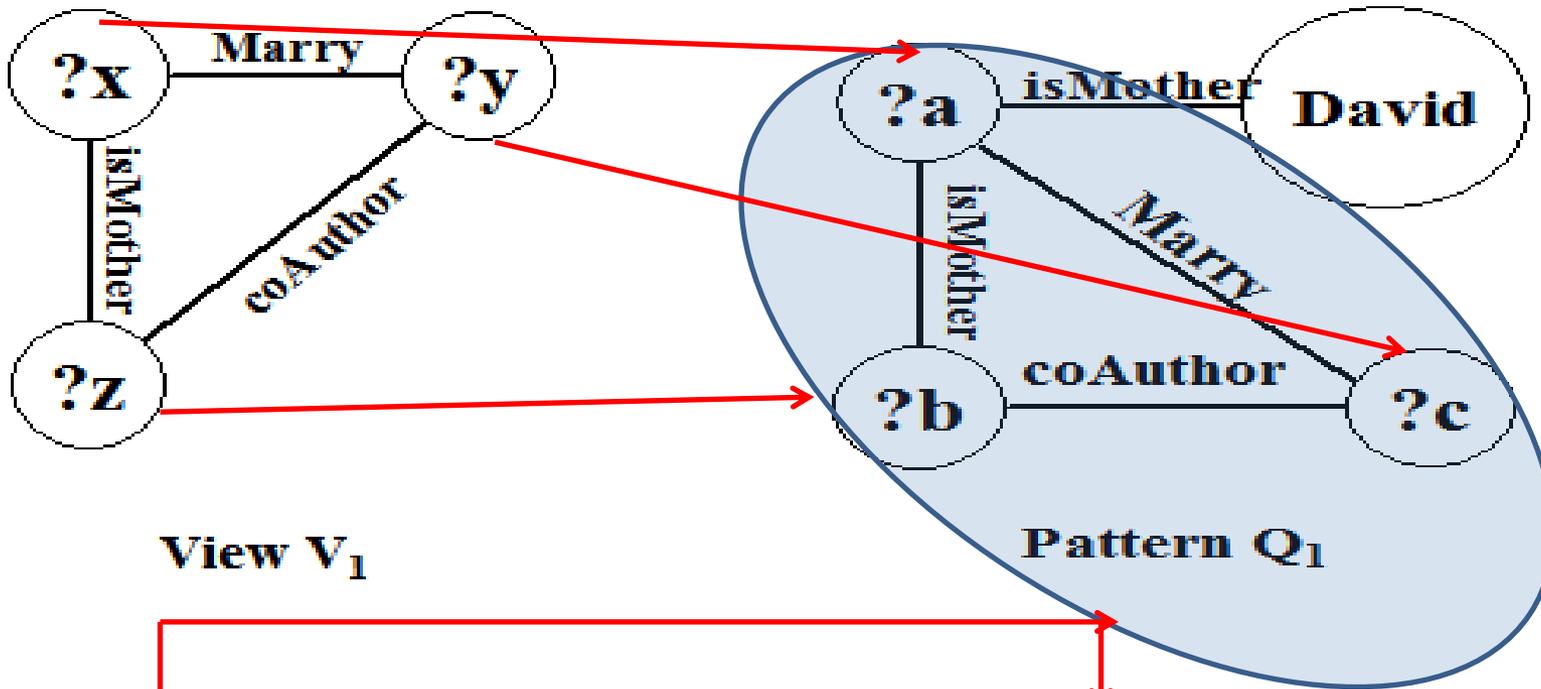
[V1]	?x	?y	?z
	Alice	Bob	Chris
	...	...	...



[Q'1]	?a	?c	?b
	Alice	Bob	Chris
	...	...	...



# RDF Pattern Rewriting using View

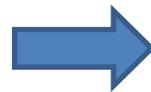


[V1]	?x	?y	?z	[Q'1]	?a	?c	?b
	Alice	Bob	Chris		Alice	Bob	Chris
	...	...	...		...	...	...



# RDF Pattern Rewriting using View

[V1]	?x	?y	?z
	Alice	Bob	Chris
	...	...	...



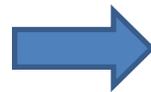
Q'1	?a	?c	?b
	Alice	Bob	Chris
	...	...	...



# RDF Pattern Rewriting using View

<b>?a</b>	<b>isMother</b>	<b>David</b>
Alice	isMother	David
....	....	....

[V1]	?x	?y	?z
	Alice	Bob	Chris
	...	...	...



<b>Q'1</b>	<b>?a</b>	<b>?c</b>	<b>?b</b>
	Alice	Bob	Chris
	...	...	...



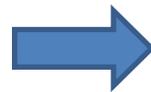
# RDF Pattern Rewriting using View

?a	isMother	David	?c	?b
Alice	isMother	David	Bob	Chris

?a	isMother	David
Alice	isMother	David
....	....	....

Join on ?a

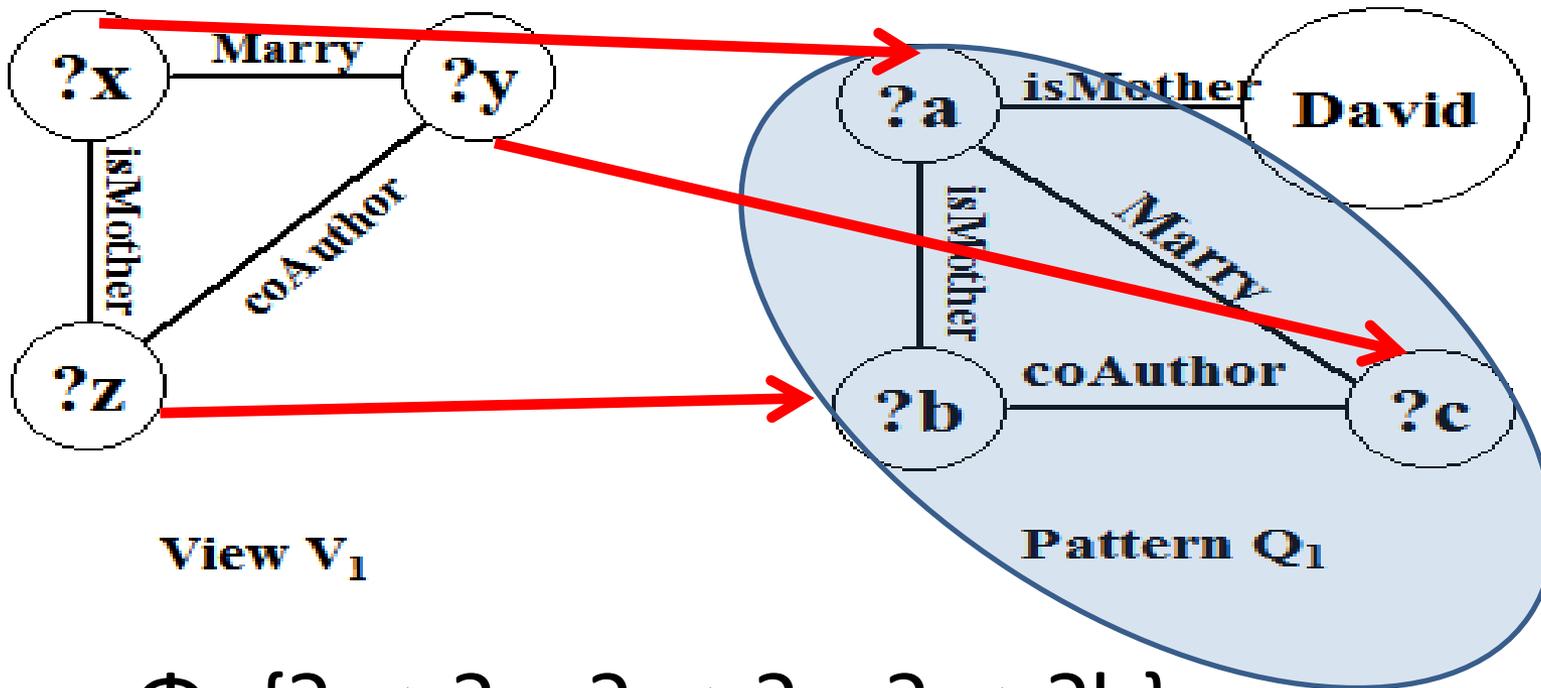
[V1]	?x	?y	?z
	Alice	Bob	Chris
	...	...	...



Q'1	?a	?c	?b
	Alice	Bob	Chris
	...	...	...



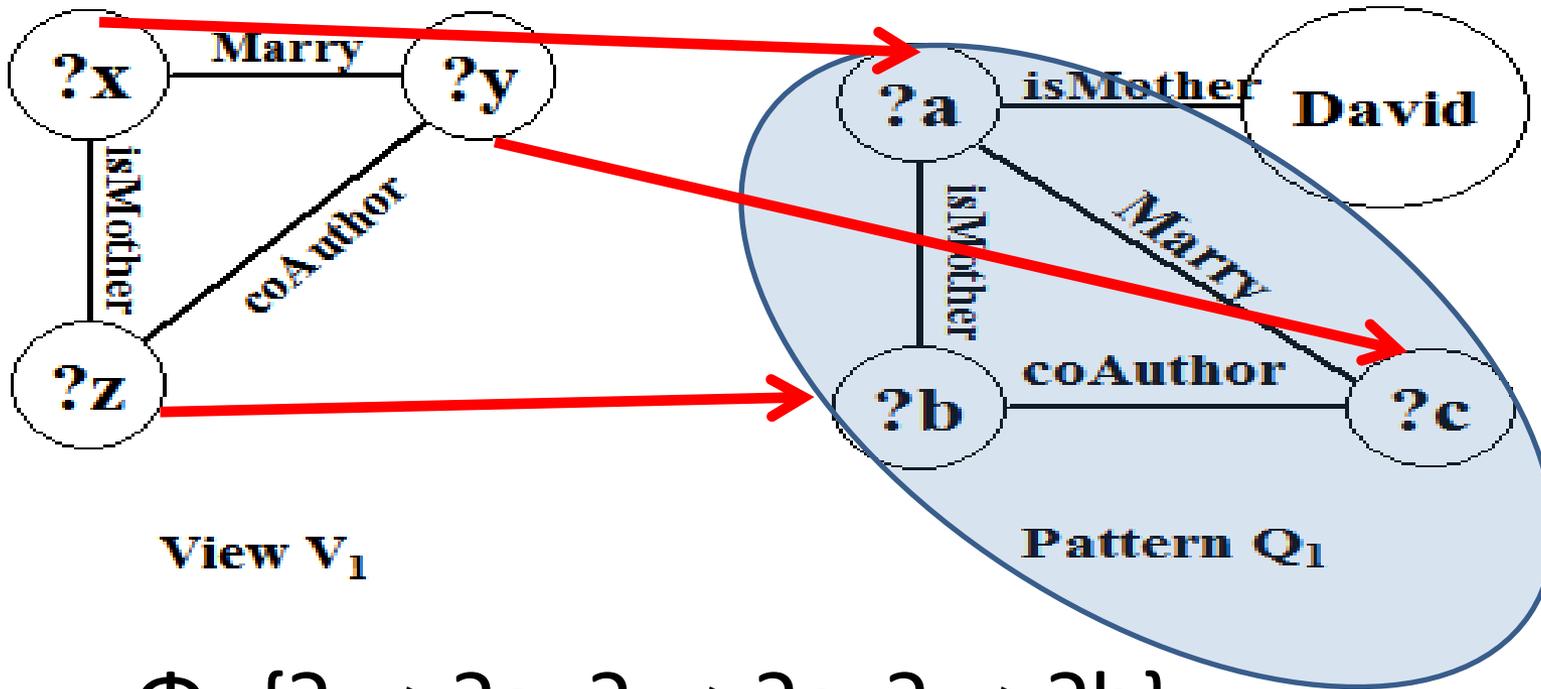
# RDF Pattern Rewriting using View



$$\Phi = \{ ?x \rightarrow ?a, ?y \rightarrow ?c, ?z \rightarrow ?b \}$$



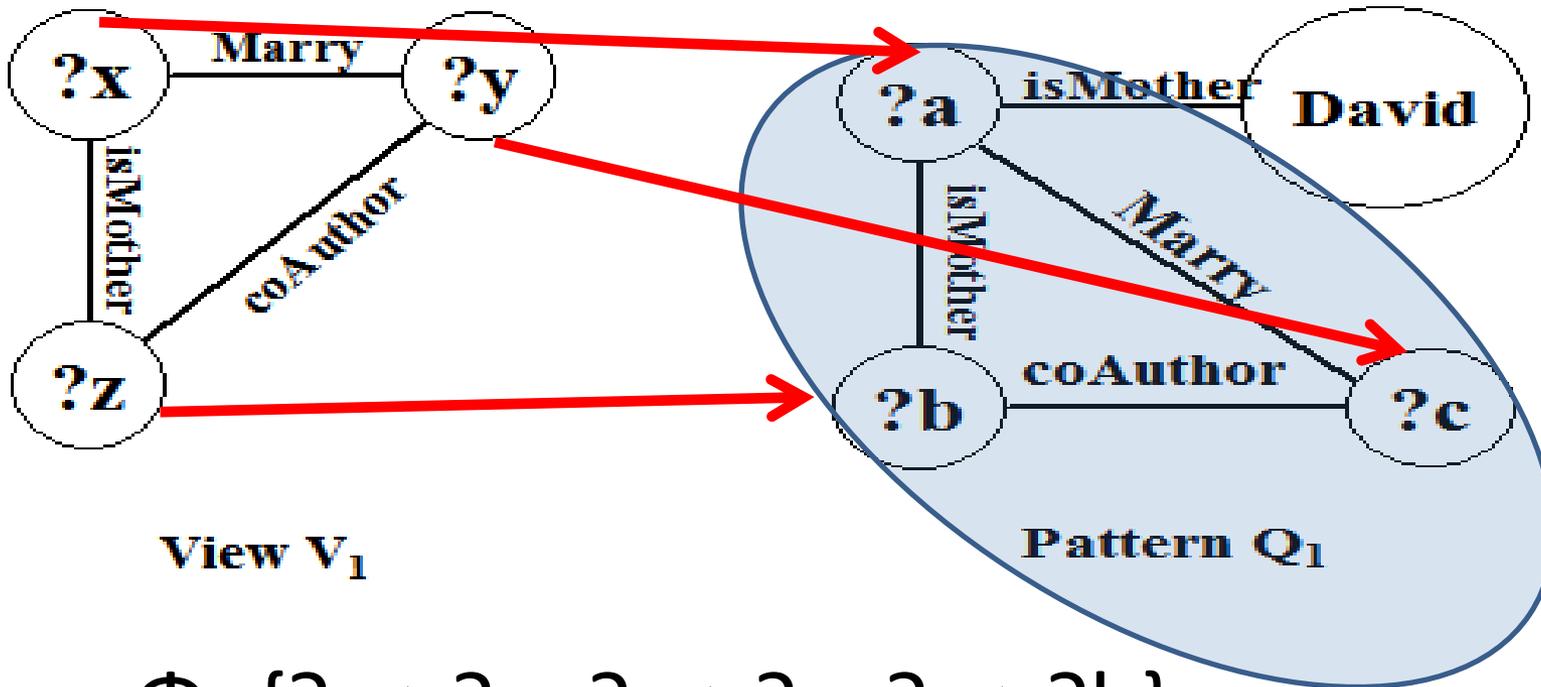
# RDF Pattern Rewriting using View



Containment mapping  $\Phi$



# RDF Pattern Rewriting using View

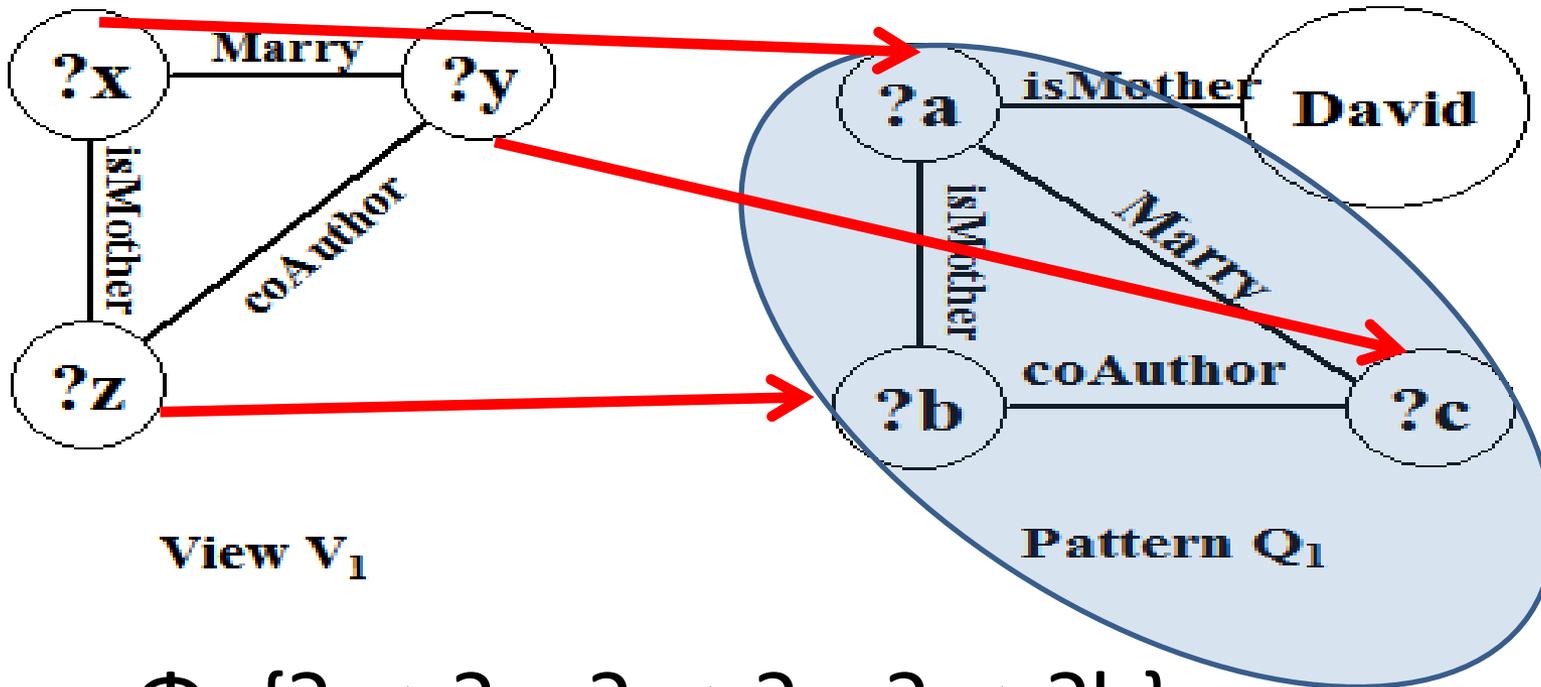


$$\Phi = \{ ?x \rightarrow ?a, ?y \rightarrow ?c, ?z \rightarrow ?b \}$$

$$Q_1 = \Phi(V_1) \text{ union } \{ (?a, \text{isMother}, \text{David}) \}$$



# RDF Pattern Rewriting using View



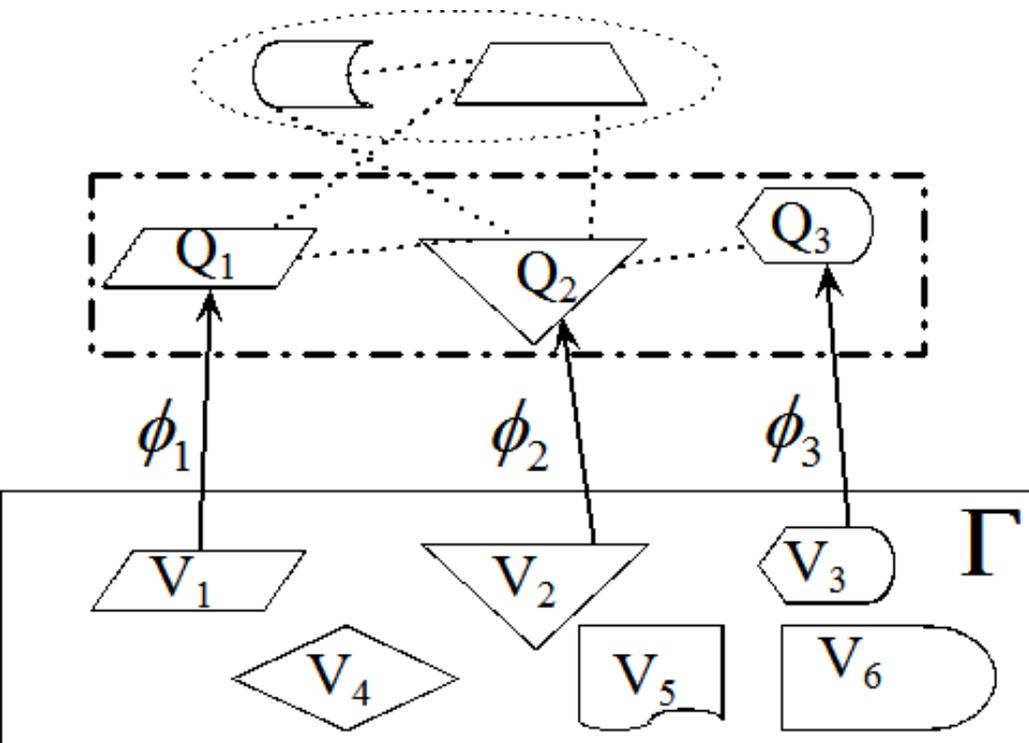
$$\Phi = \{ ?x \rightarrow ?a, ?y \rightarrow ?c, ?z \rightarrow ?b \}$$

$$Q_1 = \Phi(V_1) \text{ union } \{ (?a, \text{isMother}, \text{David}) \}$$

$$[Q_1] = \Phi([V_1]) \ominus \{ (?a, \text{isMother}, \text{David}) \}$$

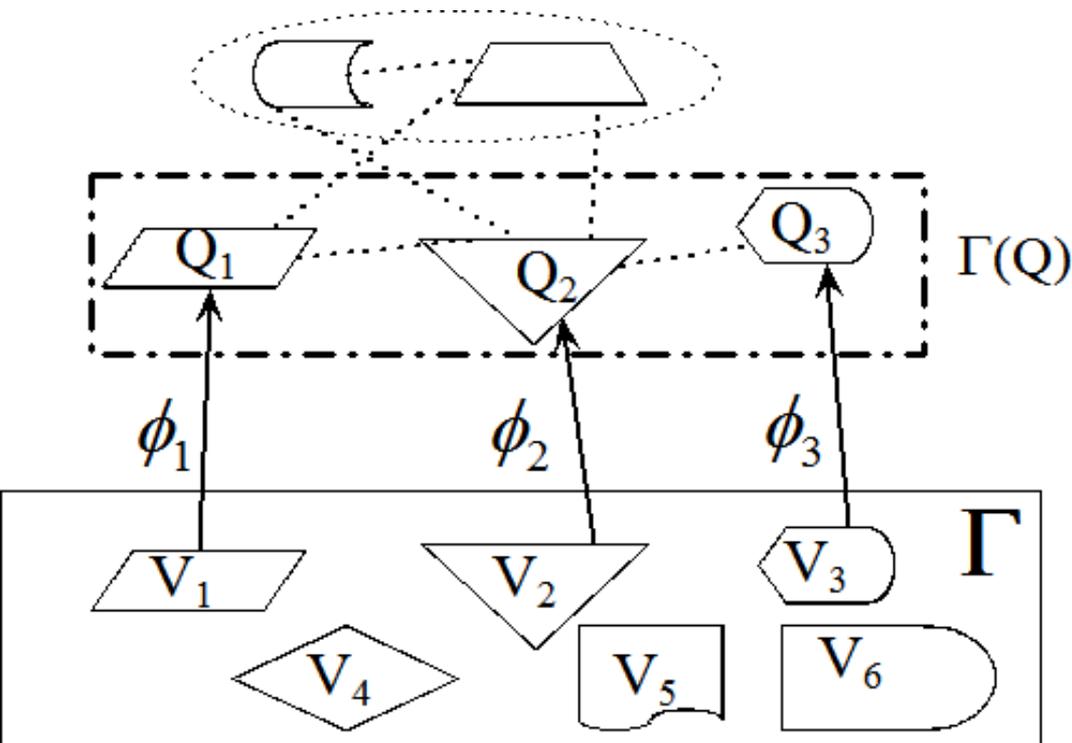


# RDF Pattern Rewriting using View



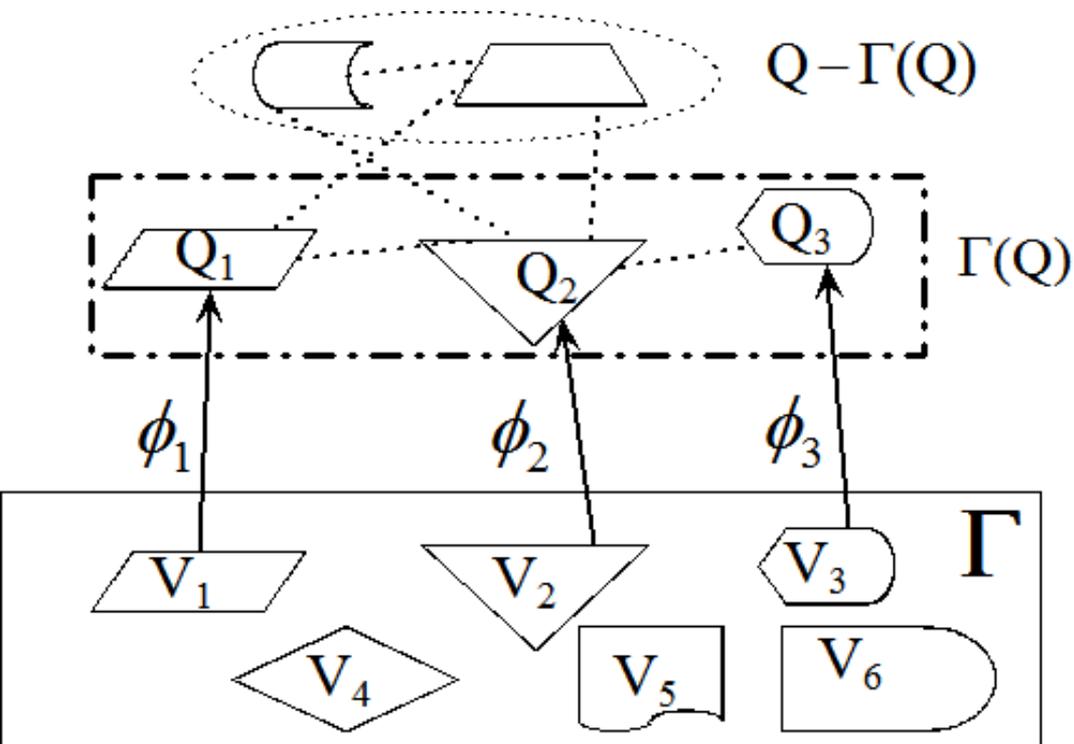


# RDF Pattern Rewriting using View



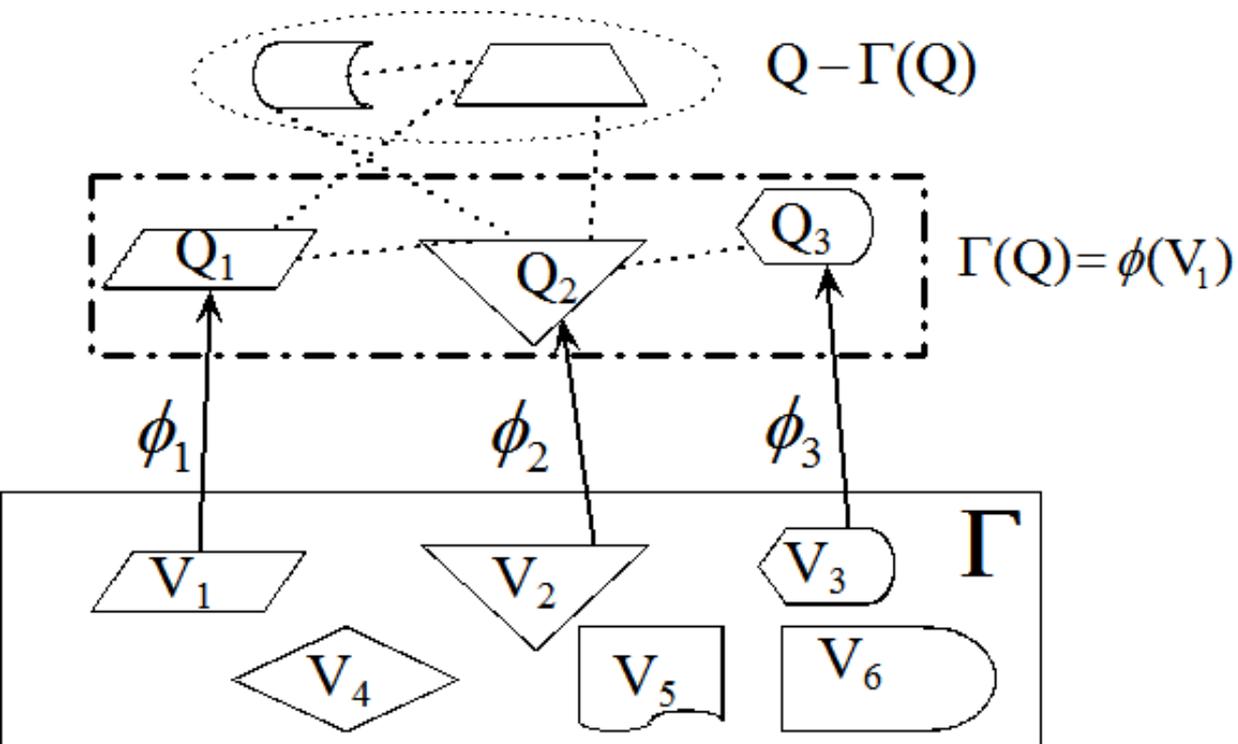


# RDF Pattern Rewriting using View



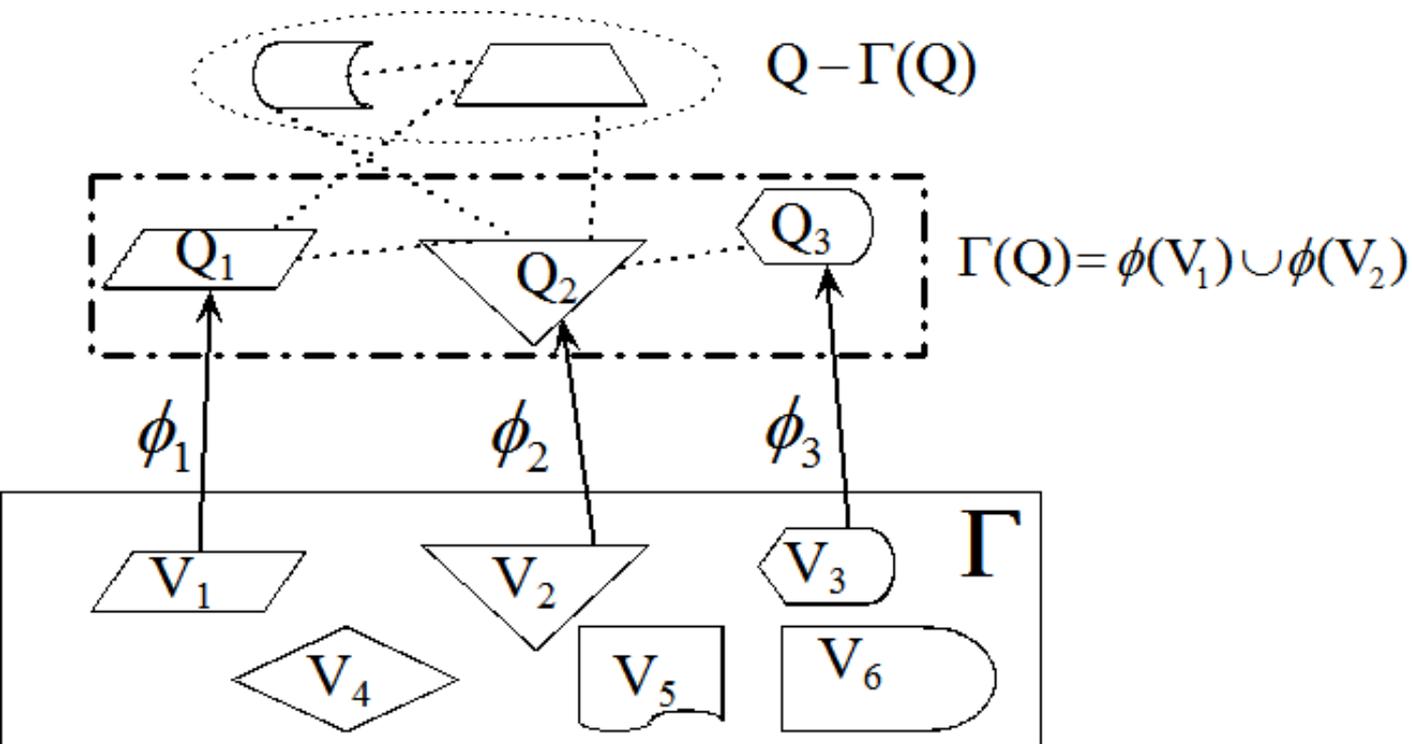


# RDF Pattern Rewriting using View



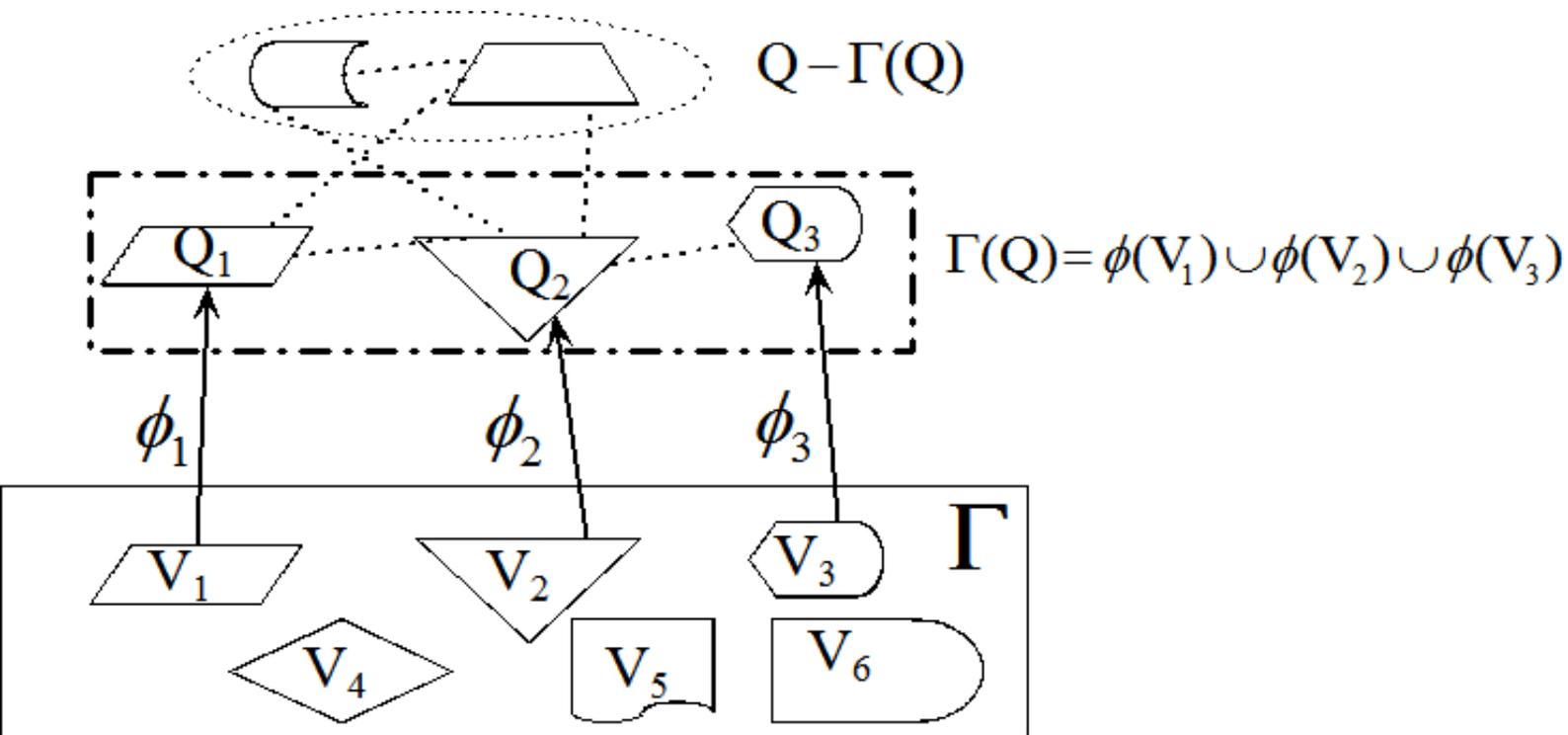


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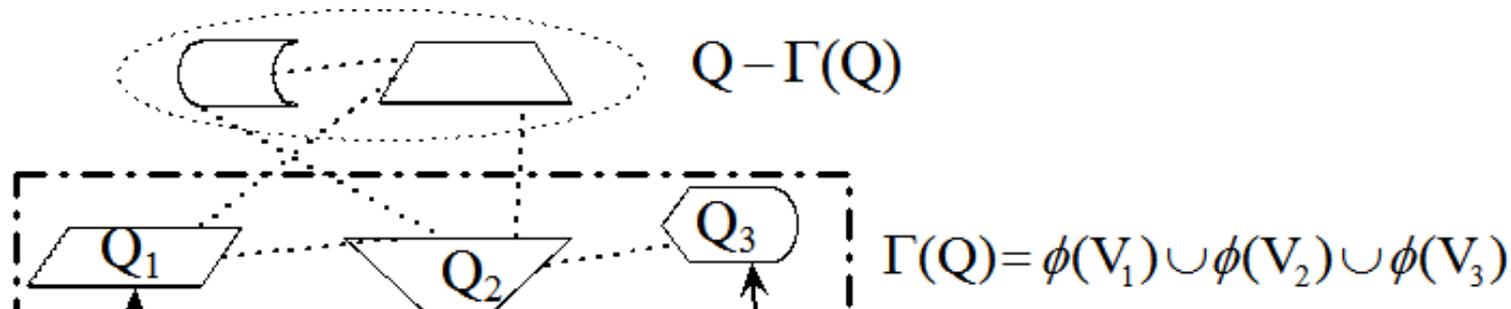


# RDF Pattern Rewriting using View





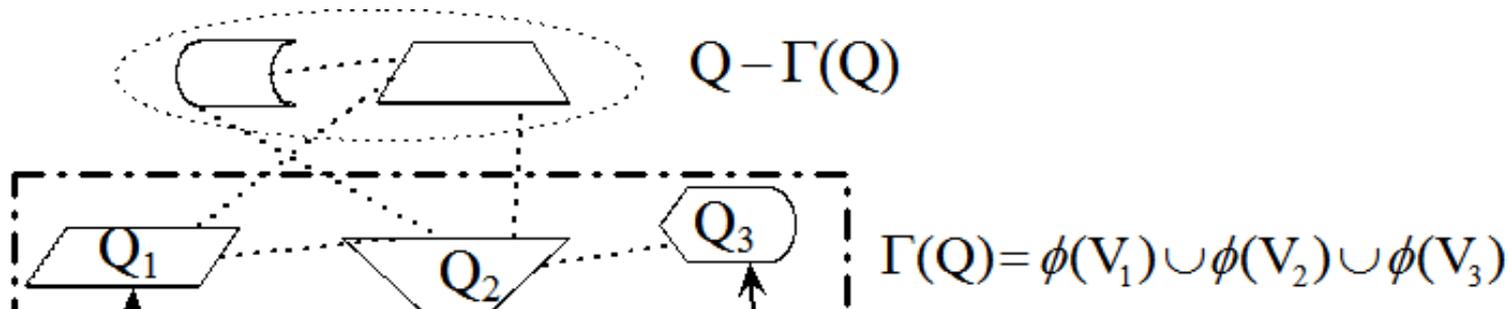
# RDF Pattern Rewriting using View



$$[Q] = [\phi_1(V_1)] \quad \times \quad [\phi_2(V_2)] \quad \times \quad [\phi_3(V_3)] \quad \times \quad [Q - \Gamma(Q)]$$



# RDF Pattern Rewriting using View



$$[Q] = [\phi_1(V_1)] \not\bowtie [\phi_2(V_2)] \not\bowtie [\phi_3(V_3)] \not\bowtie [Q - \Gamma(Q)]$$

NP-hard to find these views  $V_i$  with their  $\phi_i$ !!



# Outline

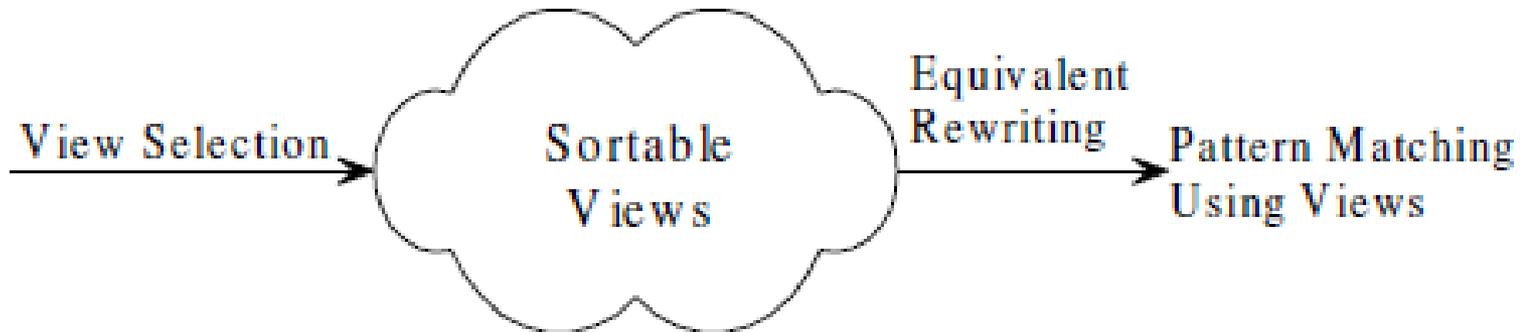
- RDF Pattern Rewriting using Views
- **Sortable View**
- Rewriting using Sortable Views
- Evaluation



# Sortable View

- The problem of rewriting pattern using views is **NP-hard** in general cases

- **Sortable View**

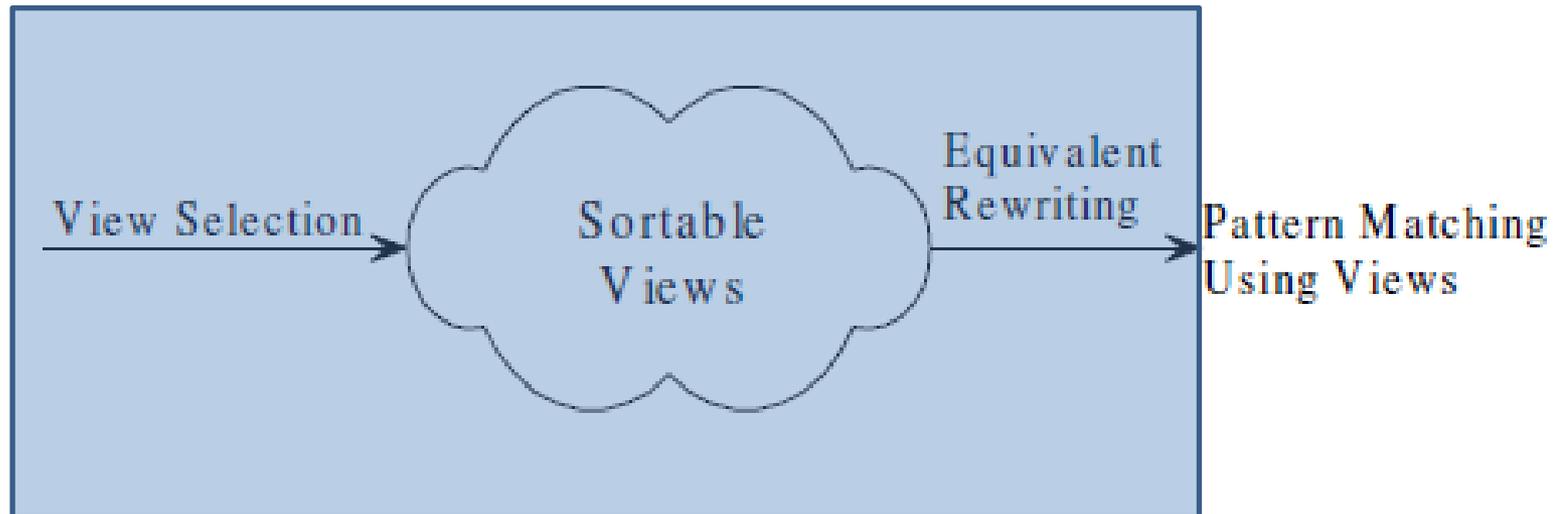




# Sortable View

- The problem of rewriting pattern using views is **NP-hard** in general cases

- **Sortable View**



# Sortable View



$V_2[1]$   
(*?p, coAuthor, ?o*)

$V_2[2]$   
(*?n, isMother, ?o*)

$V_2[3]$   
(*?m, Marry, ?n*)



# Sortable View

$V_2[1]$   
(*?p, coAuthor, ?o*)  $\curvearrowright$   $V_2[2]$   
(*?n, isMother, ?o*)  $\curvearrowright$   $V_2[3]$   
(*?m, Marry, ?n*)



# Sortable View

1. All variables are equivalent

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\curvearrowright$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\curvearrowright$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

1. All variables are equivalent
2. Two constants(strings) are ordered by their lexical order.

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\prec$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\prec$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

1. All variables are equivalent
2. Two constants(strings) are ordered by their lexical order.
3. Two triple patterns are ordered by

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\prec$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\prec$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

1. All variables are equivalent
2. Two constants(string) are ordered by their lexical order.
3. Two triple patterns are ordered by
  - Comparing their subject components first

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\prec$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\prec$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

1. All variables are equivalent
2. Two constants(string) are ordered by their lexical order.
3. Two triple patterns are ordered by
  - Comparing their subject components first
  - If equivalent, comparing their predict components

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\prec$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\prec$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

1. All variables are equivalent
2. Two constants(string) are ordered by their lexical order.
3. Two triple patterns are ordered by
  - Comparing their subject components first
  - If equivalent, comparing their predict components
  - And so on

$V_2[1]$   
 $(?p, coAuthor, ?o)$   $\prec$   $V_2[2]$   
 $(?n, isMother, ?o)$   $\prec$   $V_2[3]$   
 $(?m, Marry, ?n)$



# Sortable View

$(Frank, coAuthor, ?o') \prec (?n', isMother, ?o') \prec (Bob, Marry, ?n')$

$V_2[1] \quad \prec \quad V_2[2] \quad \prec \quad V_2[3]$   
 $(?p, coAuthor, ?o) \quad \quad \quad (?n, isMother, ?o) \quad \quad \quad (?m, Marry, ?n)$



# Sortable View

$(Frank, coAuthor, ?o') \prec (?n', isMother, ?o') \prec (Bob, Marry, ?n')$

or

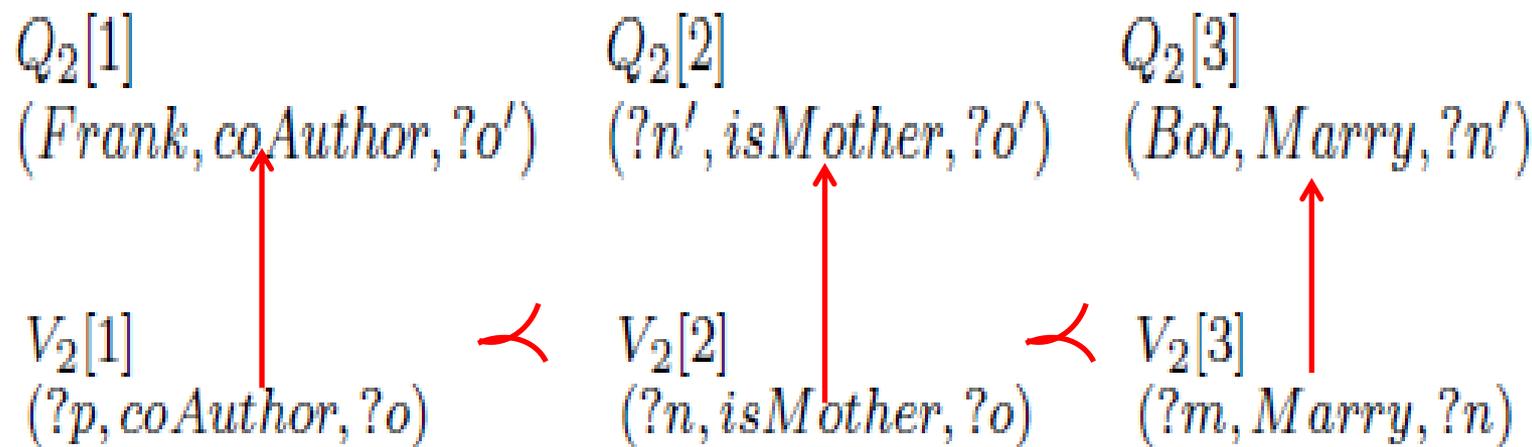
$(?n', isMother, ?o') \prec (Bob, Marry, ?n') \prec (Frank, coAuthor, ?o')$

$V_2[1] \quad \prec \quad V_2[2] \quad \prec \quad V_2[3]$   
 $(?p, coAuthor, ?o) \quad \quad \quad (?n, isMother, ?o) \quad \quad \quad (?m, Marry, ?n)$



# Sortable View

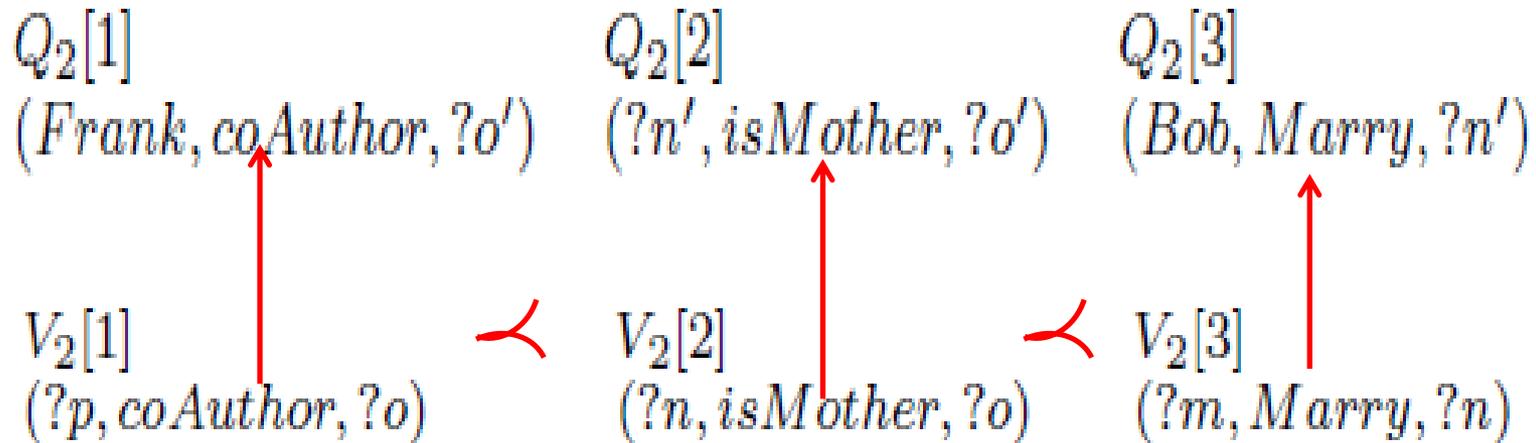
THEOREM 2. Given pattern  $Q$  against sortable view  $V$ ,  $V \sqsupseteq Q$  if and only if there is a containment mapping  $\phi$  such that  $\phi(V[i]) = \underline{Q[i]}$  for  $V$ -serialization of  $Q$ .





# Sortable View

Serializing these triple patterns in  $Q_2$  in terms of the sorted triple patterns in  $V_2$  such that  $V_2[i]$  is equivalent with  $Q_2[i]$ .  
 $Q[i]$  for  $V$ -serialization of  $Q$ .

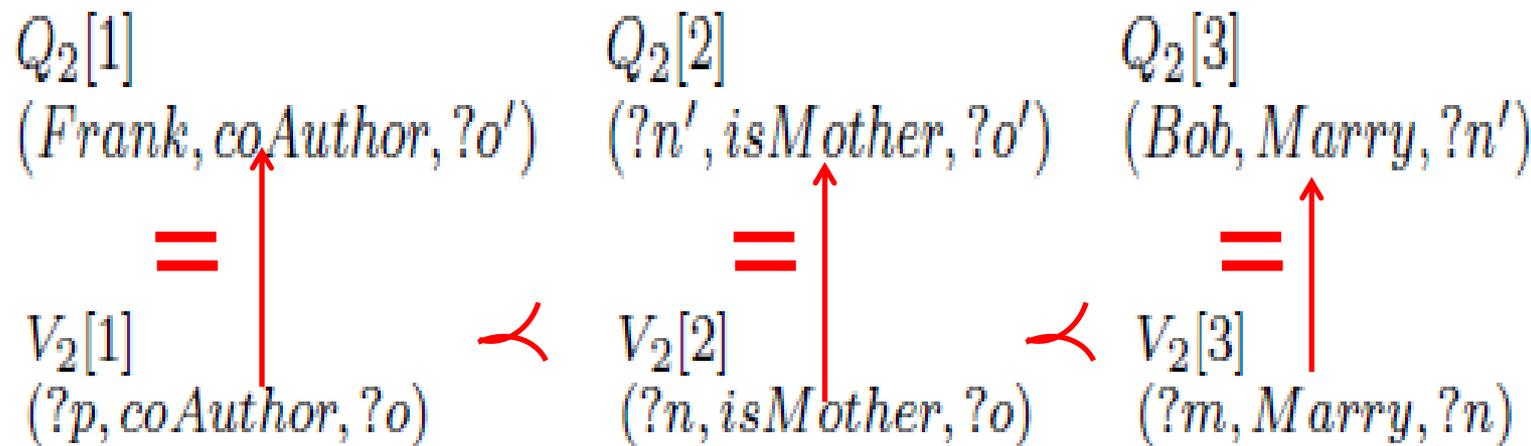




# Sortable View

Serializing these triple patterns in  $Q_2$  in terms of the sorted triple patterns in  $V_2$  such that  $V_2[i]$  is equivalent with  $Q_2[i]$ .

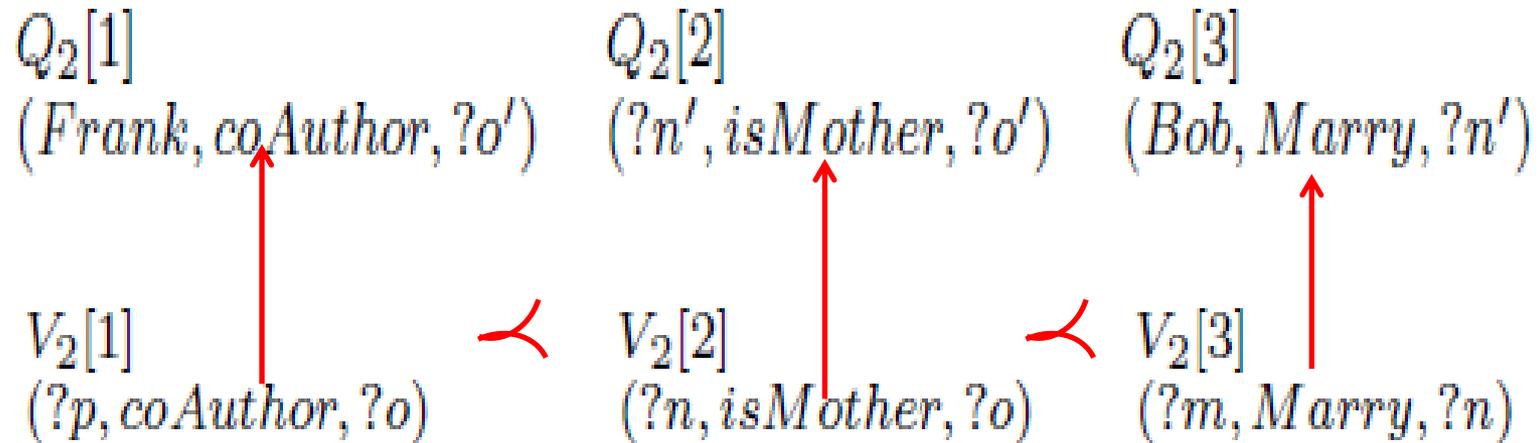
$Q_2[i]$  for  $V$ -serialization of  $Q$ .





# Sortable View

**THEOREM 2.** *Given pattern  $Q$  against sortable view  $V$ ,  $V \sqsubseteq Q$  if and only if there is a containment mapping  $\phi$  such that  $\phi(V[i]) = \underline{Q[i]}$  for  $V$ -serialization of  $Q$ .*





# Sortable View

**THEOREM 2.** *Given pattern  $Q$  against sortable view  $V$ ,  $V \sqsupseteq Q$  if and only if there is a containment mapping  $\phi$  such that  $\phi(V[i]) = \underline{Q[i]}$  for  $V$ -serialization of  $Q$ .*

**THEOREM 3.** *Pattern  $V$  is sortable if and only if there is a unique topological sort of  $\text{order}(V)$ .*

(?x, isMother, ?z)

(?y, coAuthor, ?z)

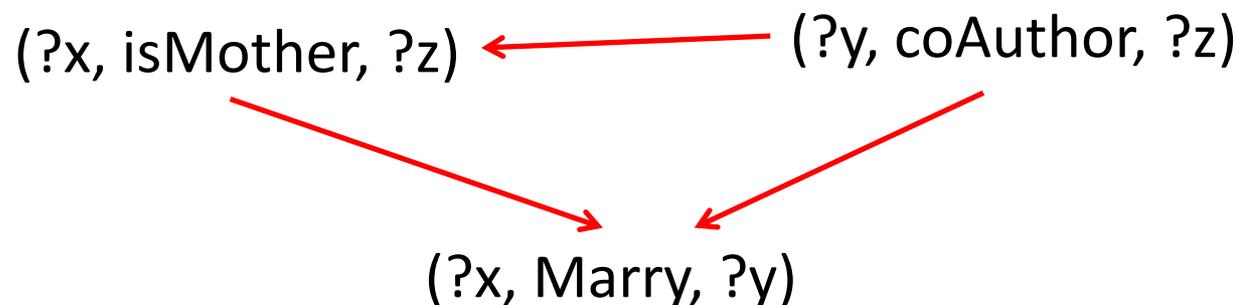
(?x, Marry, ?y)



# Sortable View

**THEOREM 2.** *Given pattern  $Q$  against sortable view  $V$ ,  $V \sqsupseteq Q$  if and only if there is a containment mapping  $\phi$  such that  $\phi(V[i]) = \underline{Q[i]}$  for  $V$ -serialization of  $Q$ .*

**THEOREM 3.** *Pattern  $V$  is sortable if and only if there is a unique topological sort of  $\text{order}(V)$ .*

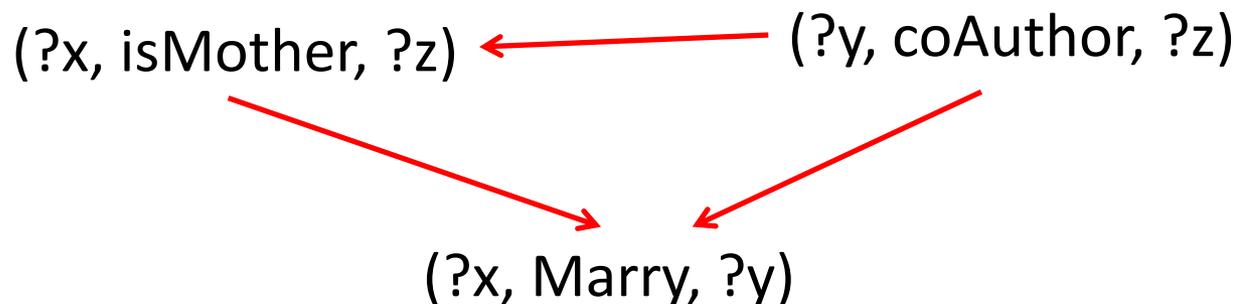




# Sortable View

**THEOREM 2.** *Given pattern  $Q$  against sortable view  $V$ ,  $V \sqsupseteq Q$  if and only if there is a containment mapping  $\phi$  such that  $\phi(V[i]) = \underline{Q[i]}$  for  $V$ -serialization of  $Q$ .*

**THEOREM 3.** *Pattern  $V$  is sortable if and only if there is a unique topological sort of  $\text{order}(V)$ .*



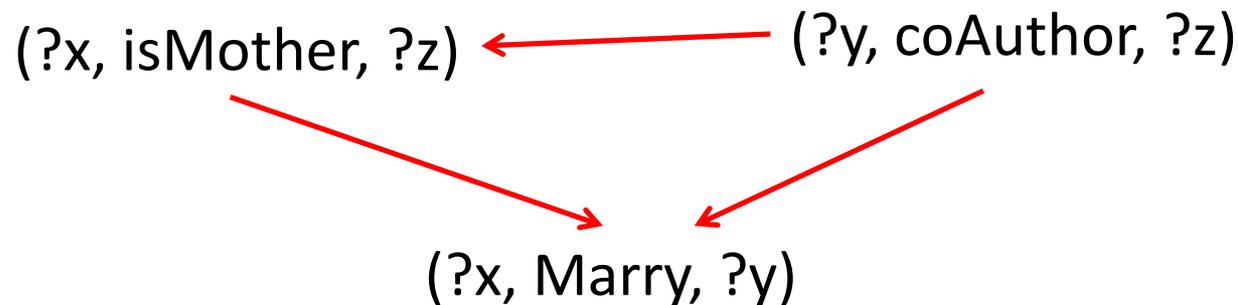
$(?y, \text{coAuthor}, ?z) < (?x, \text{isMother}, ?z) < (?x, \text{Marry}, ?y)$



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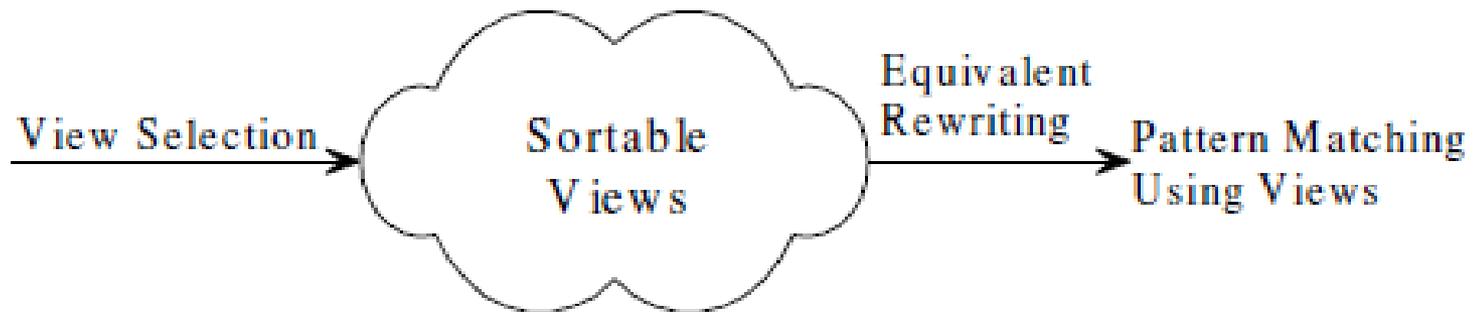
$(?y, \text{coAuthor}, ?z) < (?x, \text{isMother}, ?z) < (?x, \text{Marry}, ?y)$



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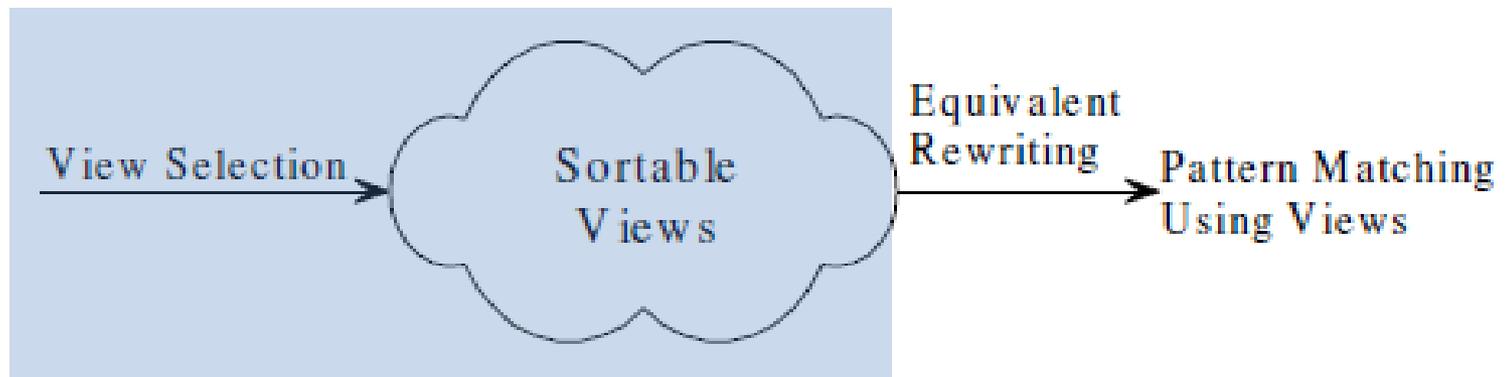




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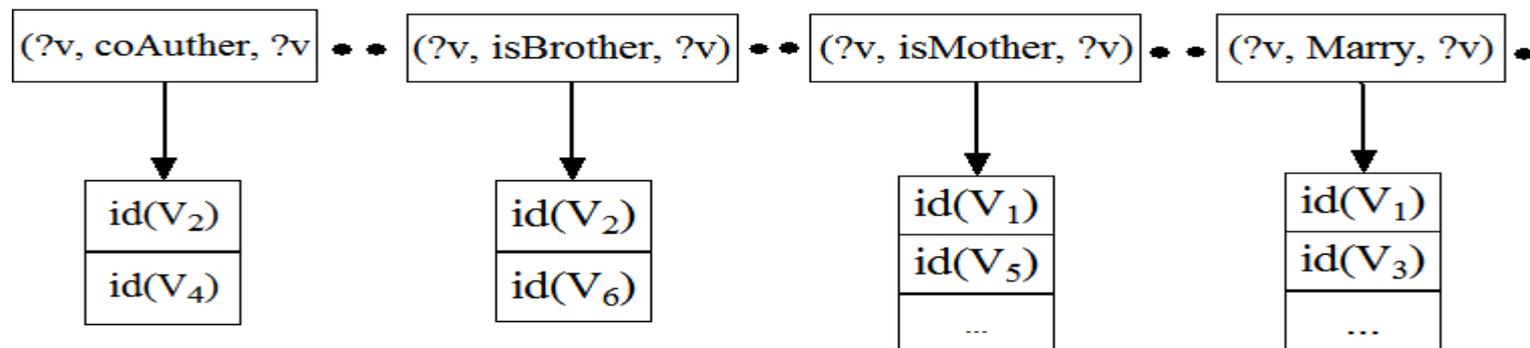


# Outline

- RDF Pattern Rewriting using Views
- Sortable View
- **Rewriting using Sortable Views**
- Evaluation

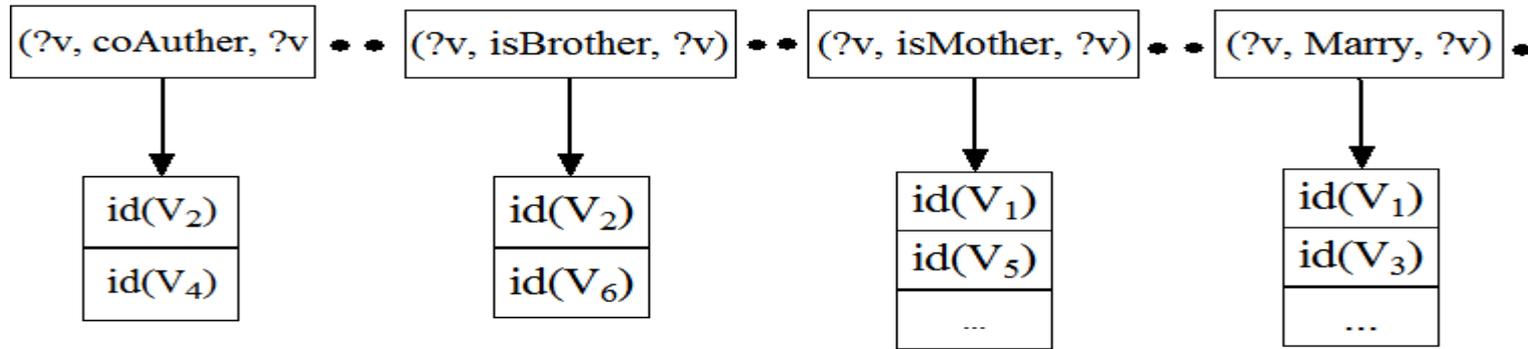


# Rewriting using Sortable Views





# Rewriting using Sortable Views



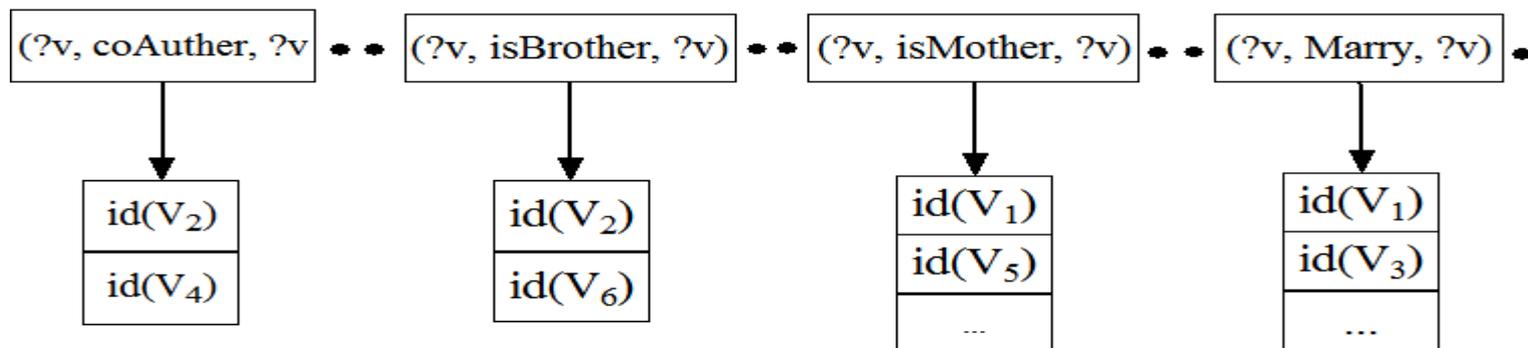
$(?c, \text{isMother}, ?d)$

$(?c, \text{Marry}, \text{Bob})$

$(?d, \text{isBrother}, ?e)$



# Rewriting using Sortable Views



(?c, isMother, ?d)

Expand each triple pattern by replacing constants with variables ?v with the **invariant:**

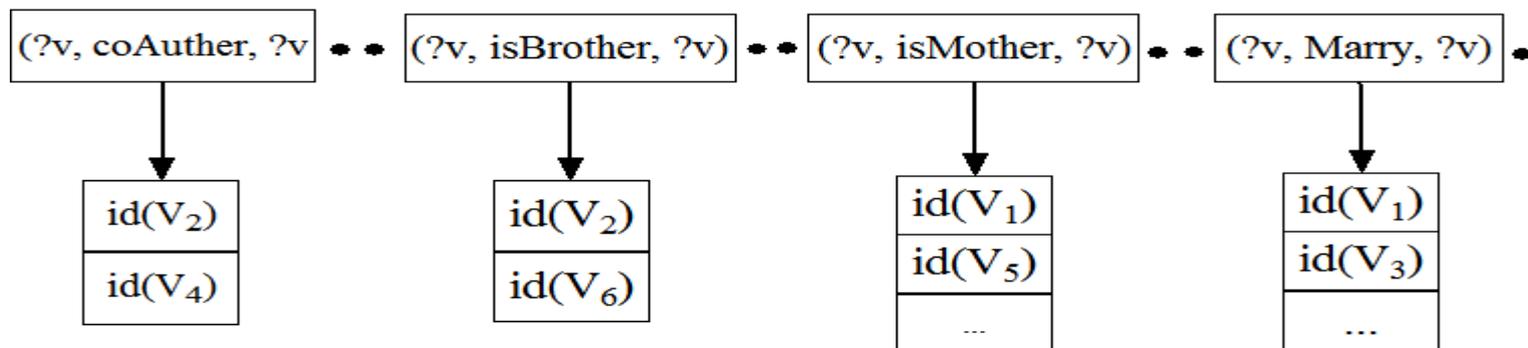
(?c, Marry, Bob)

Each triple pattern must contain at least one constant.

(?d, isBrother, ?e)



# Rewriting using Sortable Views

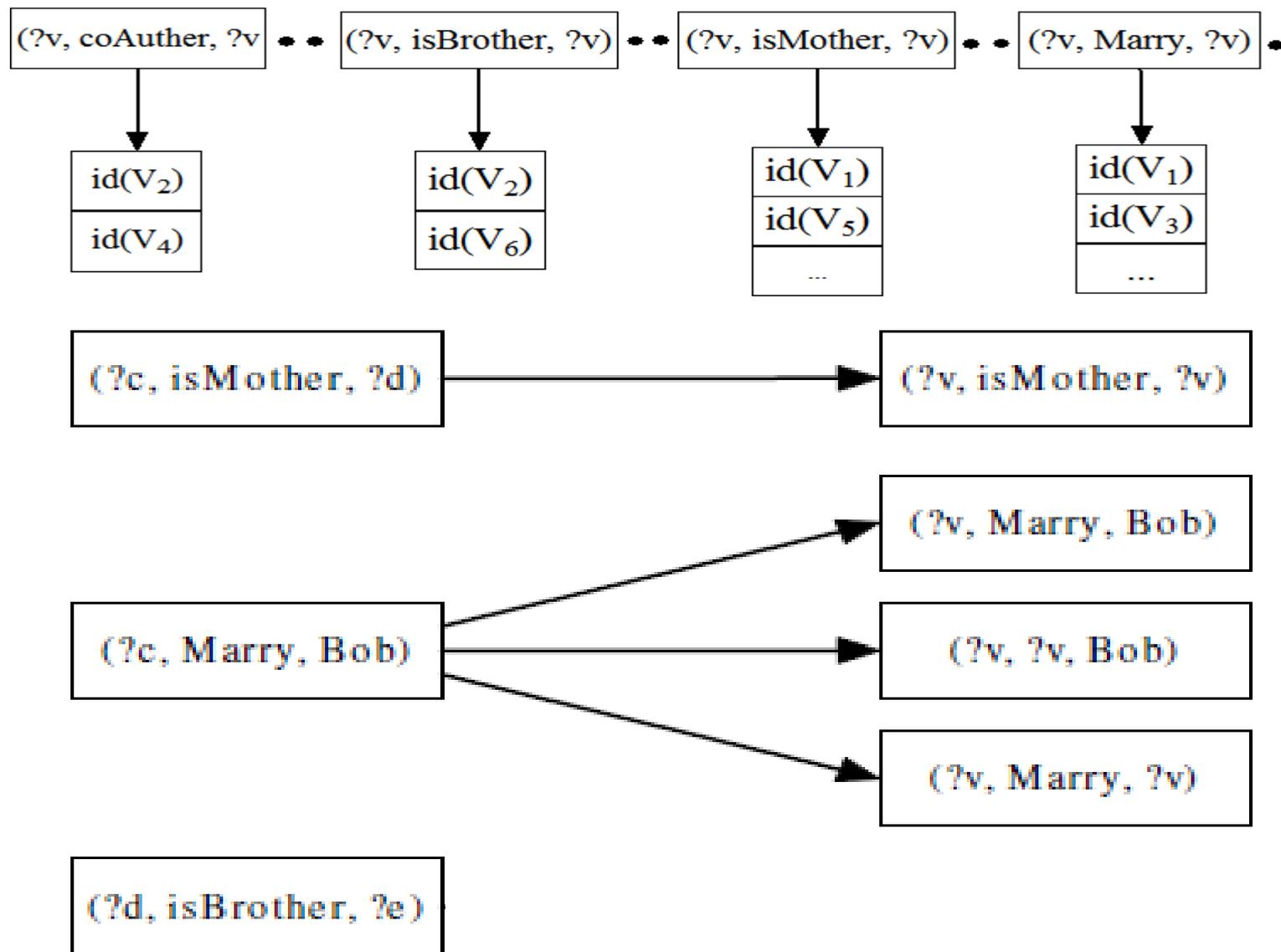


$(?c, \text{Marry}, \text{Bob})$

$(?d, \text{isBrother}, ?e)$

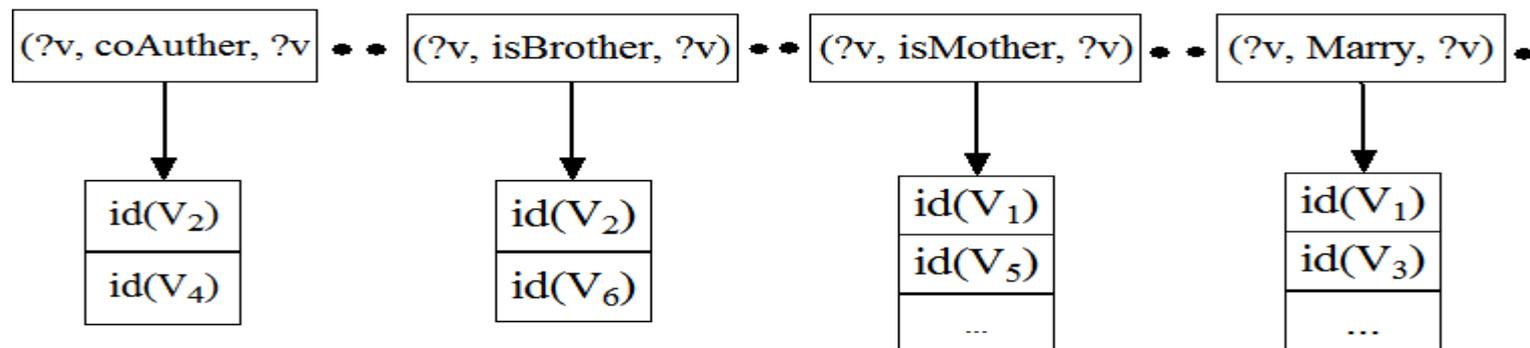


# Rewriting using Sortable Views



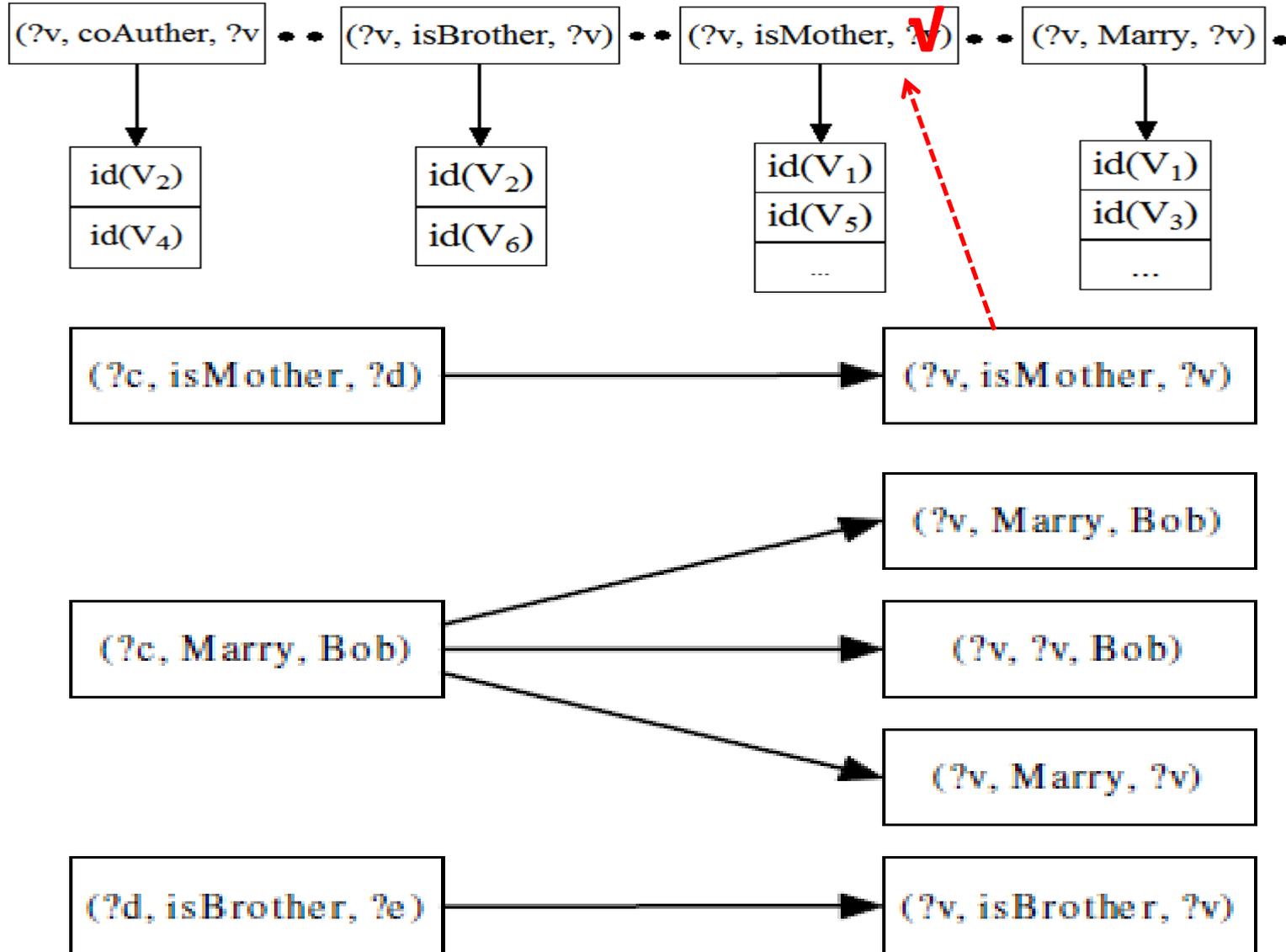


# Rewriting using Sortable Views



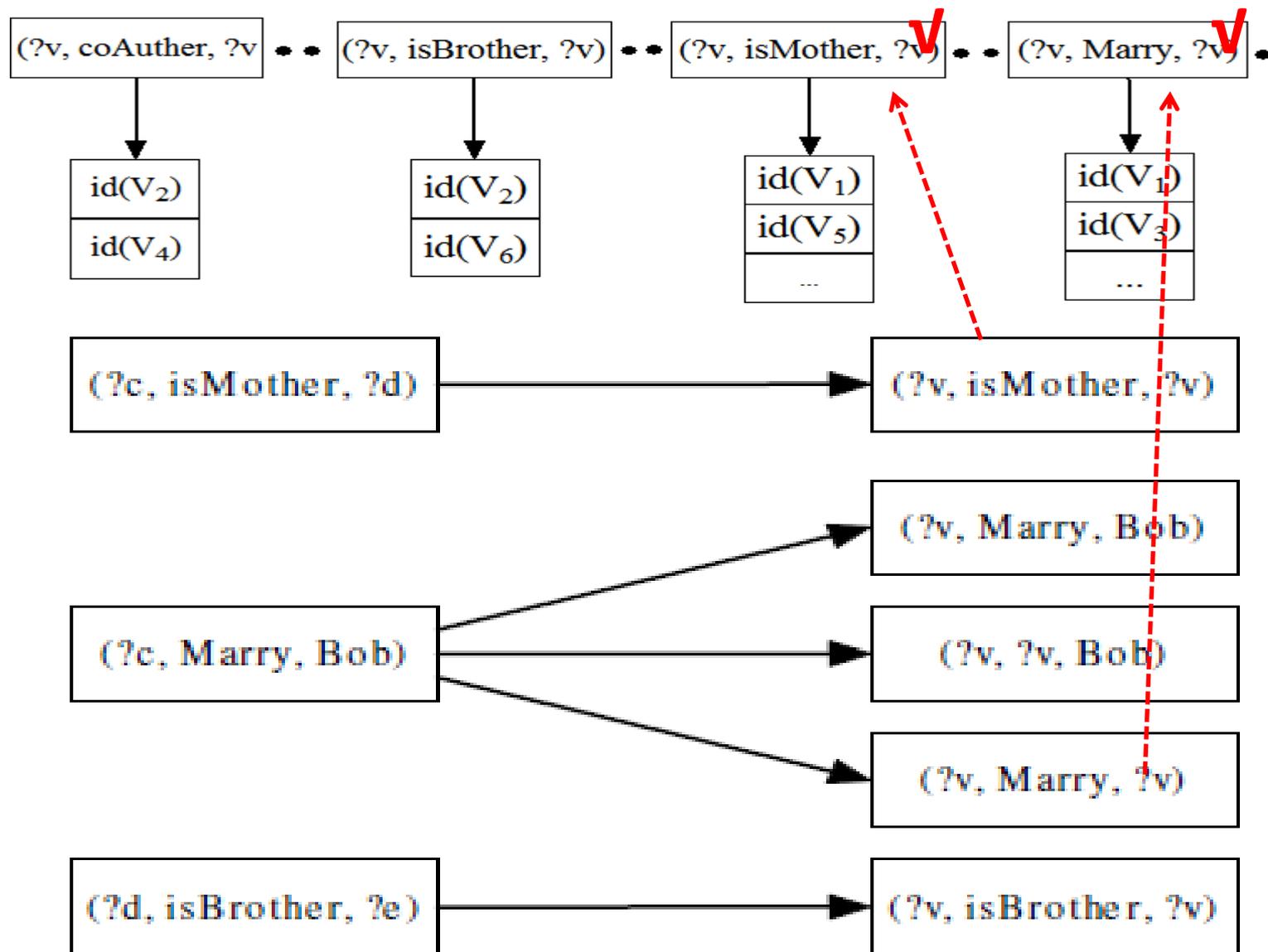


# Rewriting using Sortable Views



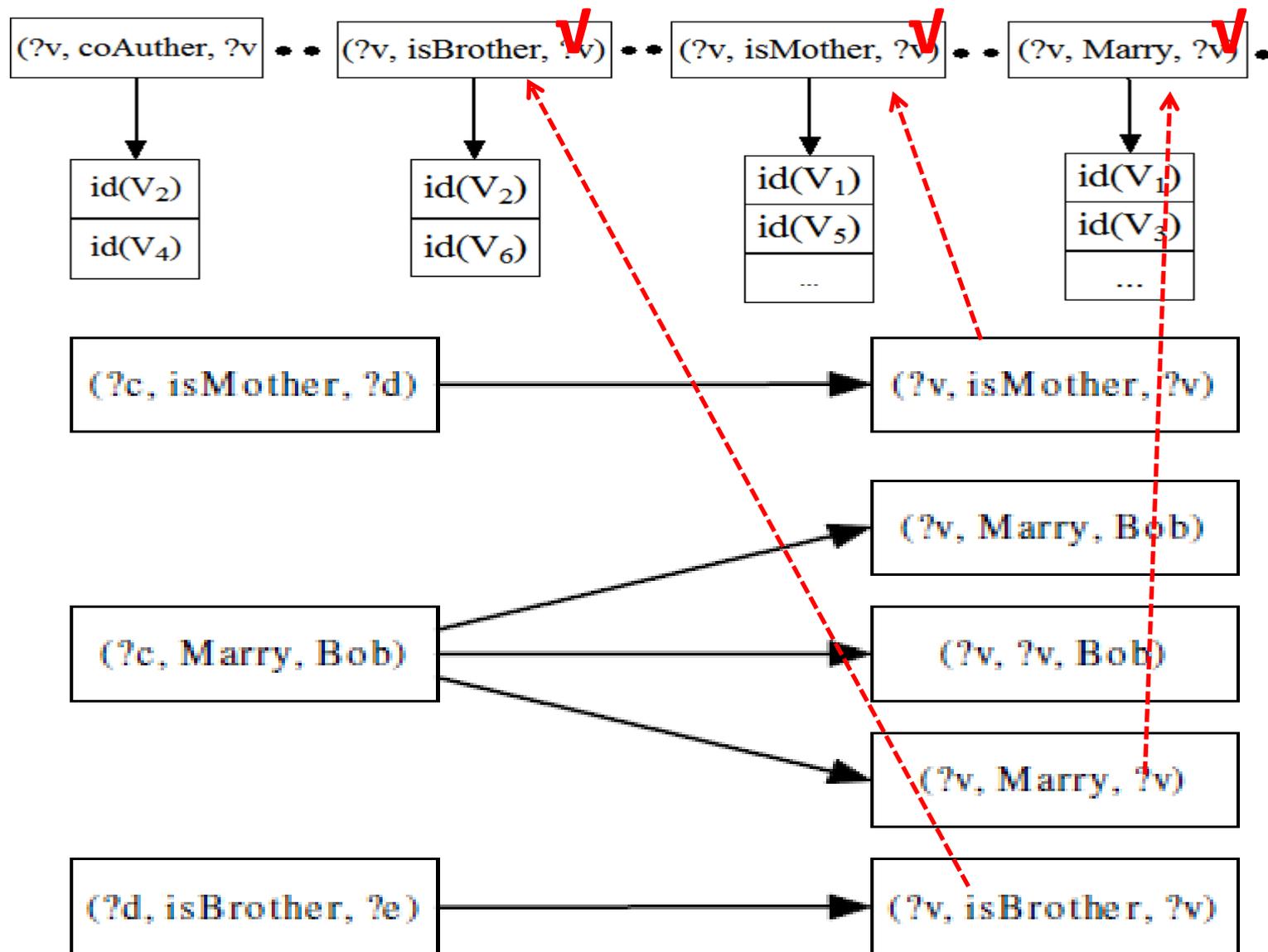


# Rewriting using Sortable Views



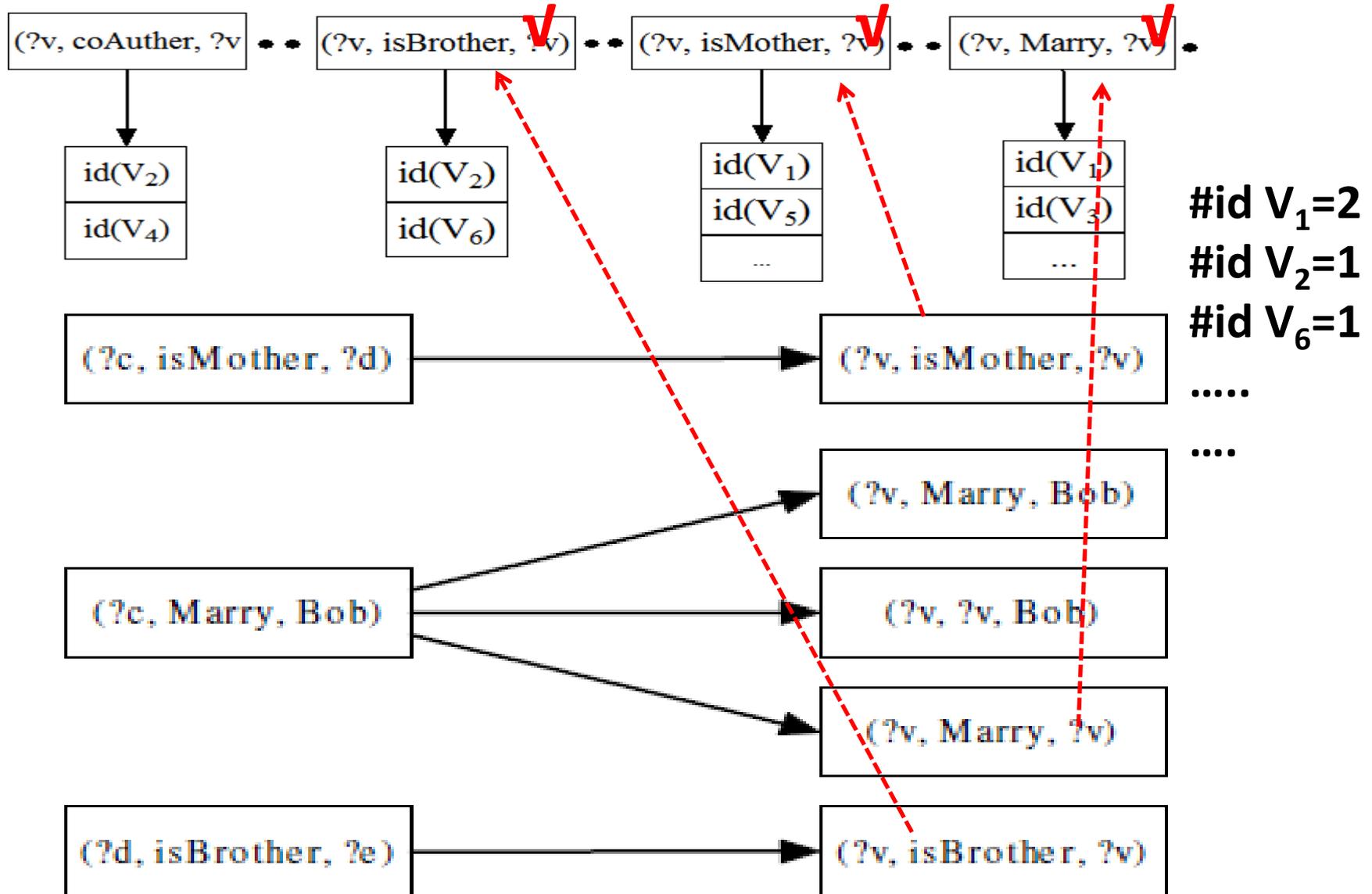


# Rewriting using Sortable Views





# Rewriting using Sortable Views





# Rewriting using Sortable Views

THEOREM 4. *Given sortable view in base  $\Gamma$ , for any  $V_i$  in  $\Gamma$ ,  $P(V_i, Q) = V_i$  if and only if  $\#id(V_i) = |V_i|$ .*

**$\#id V_1=2$**

**$\#id V_2=1$**

**$\#id V_6=1$**

.....

....



# Rewriting using Sortable Views

**THEOREM 4.** *Given sortable view in base  $\Gamma$ , for any  $V_i$  in  $\Gamma$ ,  $P(V_i, Q) = V_i$  if and only if  $\#id(V_i) = |V_i|$ .*

---

**Algorithm 4**  $rewrite^+(Q, \Gamma)$

---

- 1: for each  $q \in Q$  do
  - 2:   expand  $q$ ;
  - 3:   mark the entry equal to one of the expansions;
  - 4: end for
  - 5: count  $\#id(V_i)$  among the marked lists;
  - 6: for each  $\#id(V_i) = |V_i|$  do
  - 7:   invoke  $contain^+(V_i, Q)$  for further verification;
  - 8:   insert  $\phi_i(V_i)$  into  $\Gamma(Q)$  if  $V_i \supseteq Q$ ;
  - 9: end for
- 

**$\#id V_1=2$**

**$\#id V_2=1$**

**$\#id V_6=1$**

.....

....



# Rewriting using Sortable Views

THEOREM 4. Given sortable view in base  $\Gamma$ , for any  $V_i$  in  $\Gamma$ ,  
 $P(V_i, Q) = V_i$  if and only if  $\#id(V_i) = |V_i|$ .

$\#id V_1=2$

$\#id V_2=1$

$\#id V_6=1$

.....

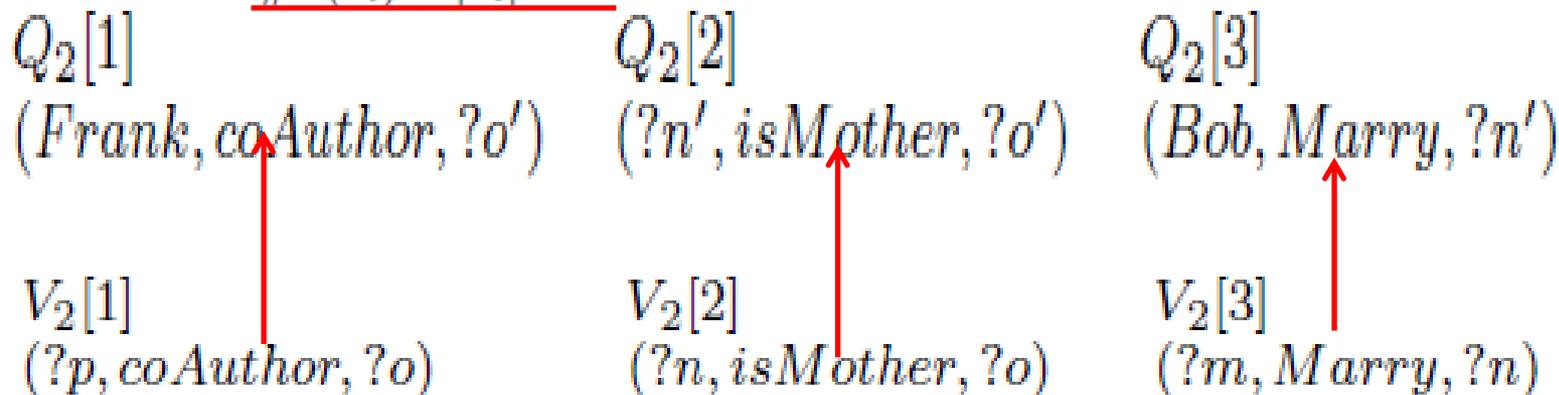
....

---

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# Rewriting using Sortable Views





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Q





# Rewriting using Sortable Views



Reduced to set cover problem with minimum cost

Q





# Outline

- RDF Pattern Rewriting using Views
- Sortable View
- Rewriting using Sortable Views
- **Evaluation**



# Evaluation

para.	description	default
$\mathcal{L}$	average length of patterns	10
$\mathcal{V}$	average number of variables	5
$M$	mode of triple generation	D
$V$	mode of variable generation	D



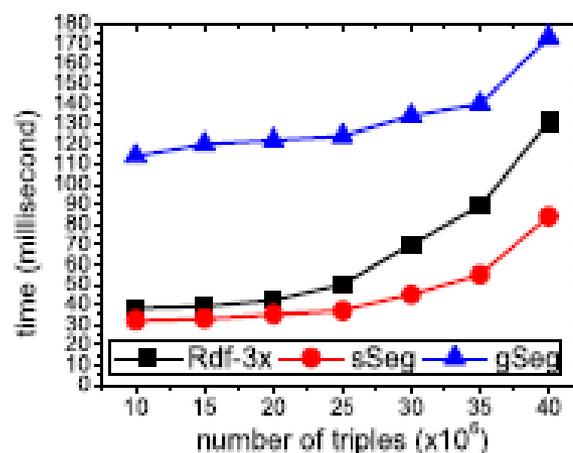
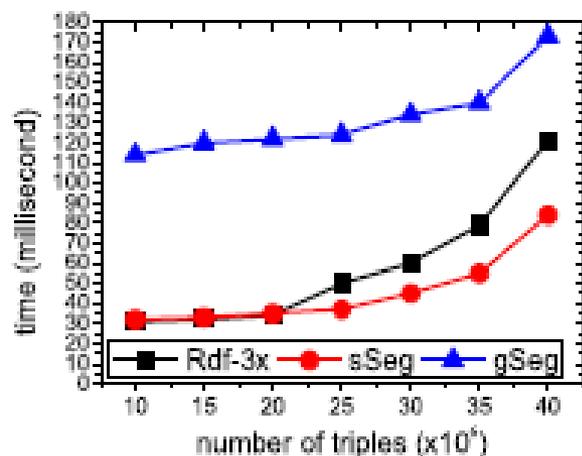
# Evaluation

para.	description	default
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$V$	mode of variable generation	D

**sSeg**  
**gSeg**  
**Rdf-3x**



# Evaluation



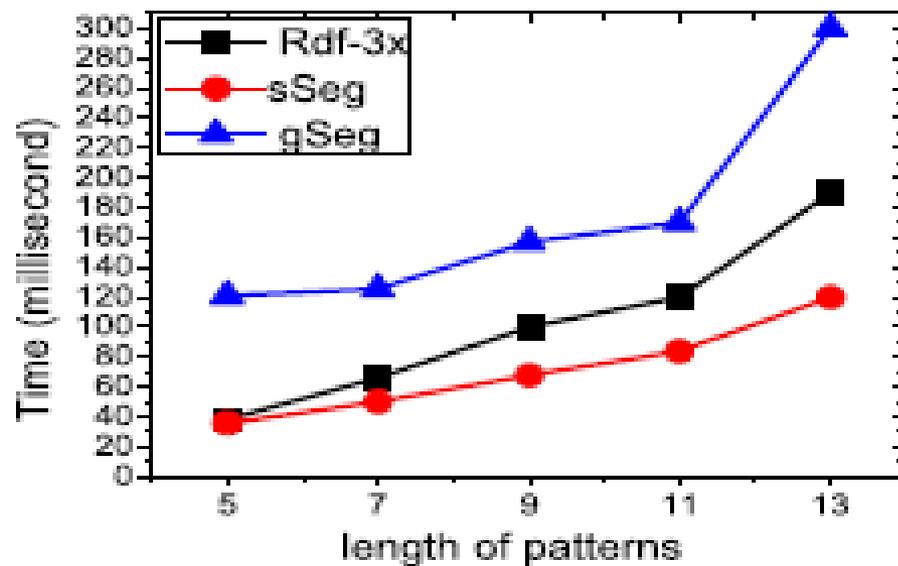
(a)  $L = 11, M=D$  and  $V=D$

(b)  $L = 11, M=F$  and  $V=F$

**Figure 5: Scalability tests on two query workloads**



# Evaluation



**Figure 8: Query length vs. time**



# Evaluation

$\mathcal{L}$ on $\Gamma_g$	$\mathcal{L}$ on $\Gamma_s$	$C_{\Gamma_s}$	sSeg	gSeg
3	3	30%	1.1	110.6
5	4	33%	2.1	121
7	6	37%	3.4	204
9	7.5	41%	3.6	417

**Table 4: Average length  $L$  of patterns in  $\Gamma_g$  vs. Time of pattern segmentation (milliseconds,  $C_{\Gamma_g} = 30\%$ )**



# Outline

- RDF Pattern Rewriting using Views
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# Conclusion

- **RDF Pattern Rewriting using Sortable Views**
  - Simplify containment mapping
  - Using inverted index structure to speed up rewriting
- **Optimum rewriting is also considered**
- **Evaluation**



**Thanks!**  
Thanks!  
**Q&A**  
Q&A