Digital Search Trees & Binary Tries

- Analog of radix sort to searching.
- Keys are binary bit strings.
  - Fixed length – 0110, 0010, 1010, 1011.
  - Variable length – 01, 00, 101, 1011.
  - IPv4 – 32 bit IP address.
  - IPv6 – 128 bit IP address.
Digital Search Tree

- Assume fixed number of bits.
- Not empty
  - Root contains one dictionary pair (any pair).
  - All remaining pairs whose key begins with a 0 are in the left subtree.
  - All remaining pairs whose key begins with a 1 are in the right subtree.
  - Left and right subtrees are digital subtrees on remaining bits.
Example

- Start with an empty digital search tree and insert a pair whose key is 0110.

- Now, insert a pair whose key is 0010.
Example

• Now, insert a pair whose key is 1001.
Example

- Now, insert a pair whose key is 1011.
Example

• Now, insert a pair whose key is 0000.
Search/Insert/Delete

• Complexity of each operation is $O(\#\text{bits in a key})$.
• $\#\text{key comparisons} = O(\text{height})$.
• Expensive when keys are very long.
Binary Trie

• Information Retrieval.
• At most one key comparison per operation.
• Fixed length keys.
  ▪ Branch nodes.
    • Left and right child pointers.
    • No data field(s).
  ▪ Element nodes.
    • No child pointers.
    • Data field to hold dictionary pair.
Example

At most one key comparison for a search.
Variable Key Length

• Left and right child fields.
• Left and right pair fields.
  ▪ Left pair is pair whose key terminates at root of left subtree or the single pair that might otherwise be in the left subtree.
  ▪ Right pair is pair whose key terminates at root of right subtree or the single pair that might otherwise be in the right subtree.
  ▪ Field is null otherwise.
Example

At most one key comparison for a search.
Fixed Length Insert

Insert 0111.  Zero compares.
Fixed Length Insert

Insert 1101.
Fixed Length Insert

Insert 1101.
Fixed Length Insert

Insert 1101.

One compare.
Delete 0111.
Delete 0111. One compare.
Delete 1100.
Delete 1100.
Delete 1100.
Delete 1100.
Fixed Length Delete

Delete 1100.

One compare.
Compressed Binary Tries

- No branch node whose degree is 1.
- Add a bit# field to each branch node.
- bit# tells you which bit of the key to use to decide whether to move to the left or right subtrie.
Binary Trie

bit# field shown in black outside branch node.
Compressed Binary Trie

The bit# field shown in black outside branch node.
#branch nodes = n – 1.
Insert 0010.
Insert 0100.
Insert

```
0 1 0001 1000 1001 0 0 0 0 0100 1 2 0 1 1
3 0 1 0 1 0001 0010 0011 4
```

```
1 1100 1101 2 0 0 0 0 1 1 4
4 1000 1001 1100
```
Delete

Delete 0010.
Delete 1001.
Delete
Higher Order Tries

- Key = Social Security Number.
  - 441-12-1135
  - 9 decimal digits.
- 10-way trie (order 10 trie).

Height <= 10.
Social Security Trie

• **10-way trie**
  - Height $\leq 10$.
  - Search $\Rightarrow \leq 9$ branches on digits plus 1 compare.

• **100-way trie**
  - 441-12-1135
  - Height $\leq 6$.
  - Search $\Rightarrow \leq 5$ branches on digits plus 1 compare.
Social Security AVL & Red-Black

- Red-black tree
  - Height $\leq 2\log_2 10^9 \sim 60$.
  - Search $\Rightarrow \leq 60$ compares of 9 digit numbers.

- AVL tree
  - Height $\leq 1.44\log_2 10^9 \sim 40$.
  - Search $\Rightarrow \leq 40$ compares of 9 digit numbers.

- Best binary tree.
  - Height $= \log_2 10^9 \sim 30$. 
Compressed Social Security Trie

Branch Node Structure

- **char#** = character/digit used for branching.
  - Equivalent to **bit#** field of compressed binary trie.
- **#ptr** = # of nonnull pointers in the node.
Insert

Insert 012345678.

Insert 015234567.

Null pointer fields not shown.
Insert 015231671.

Insert 015234567.
Insert 079864231.
Insert 012345618.
Insert 011917352.
Delete 011917352.
Delete 012345678.
Delete 015231671.

Delete 079864231.
Variable Length Keys

Problem arises only when one key is a (proper) prefix of another.
Variable Length Keys

Add a special end of key character (#) to each key to eliminate this problem.

Insert 0123.
Variable Length Keys

Insert 0123.

End of key character (#) not shown.
Tries With Edge Information

• Add a new field (element) to each branch node.
• New field points to any one of the element nodes in the subtree.
• Use this pointer on way down to figure out skipped-over characters.
Example

element field shown in blue.